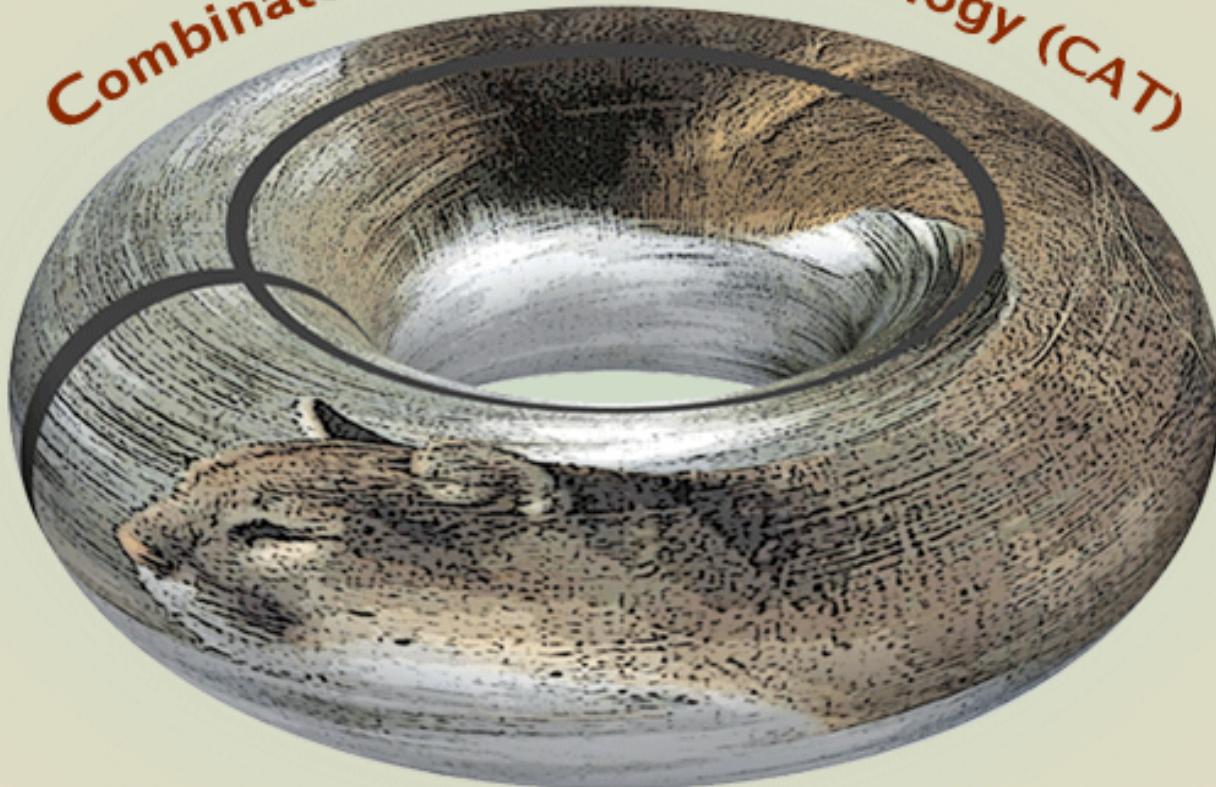
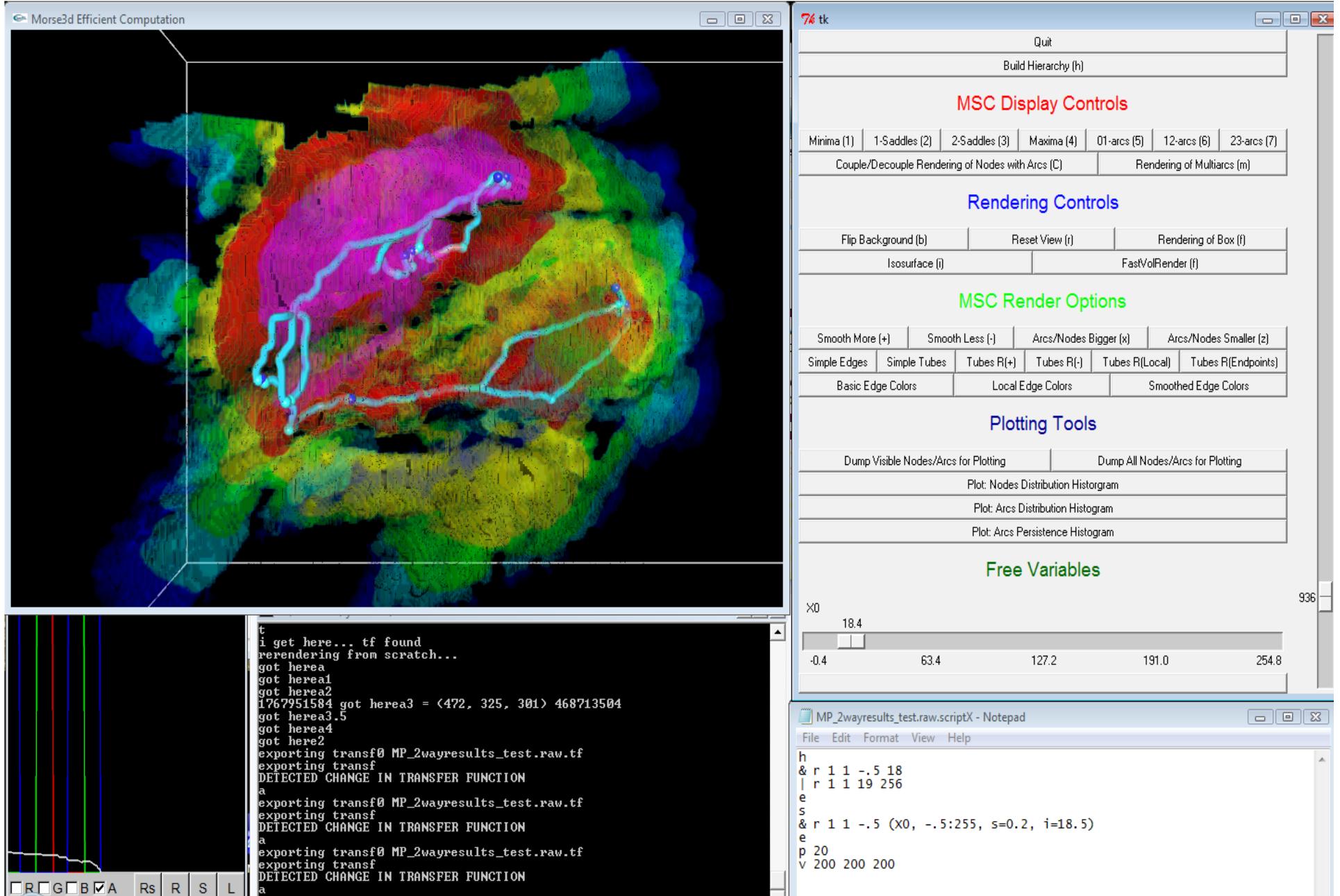


Combinatorial Algebraic Topology (CAT)

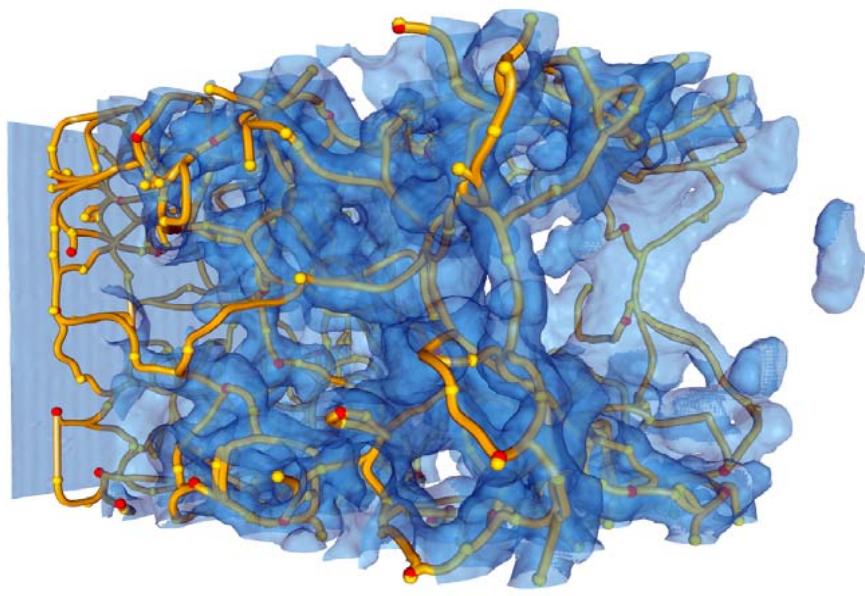
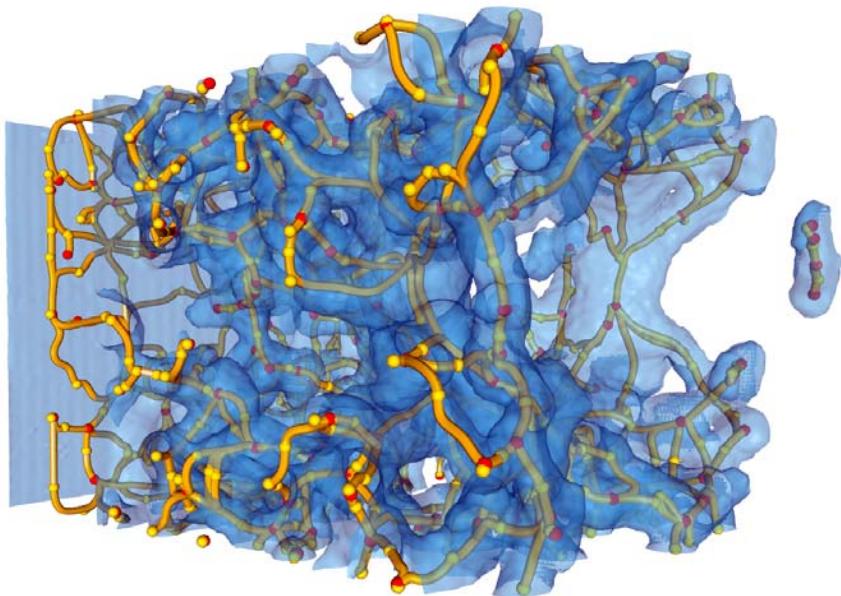
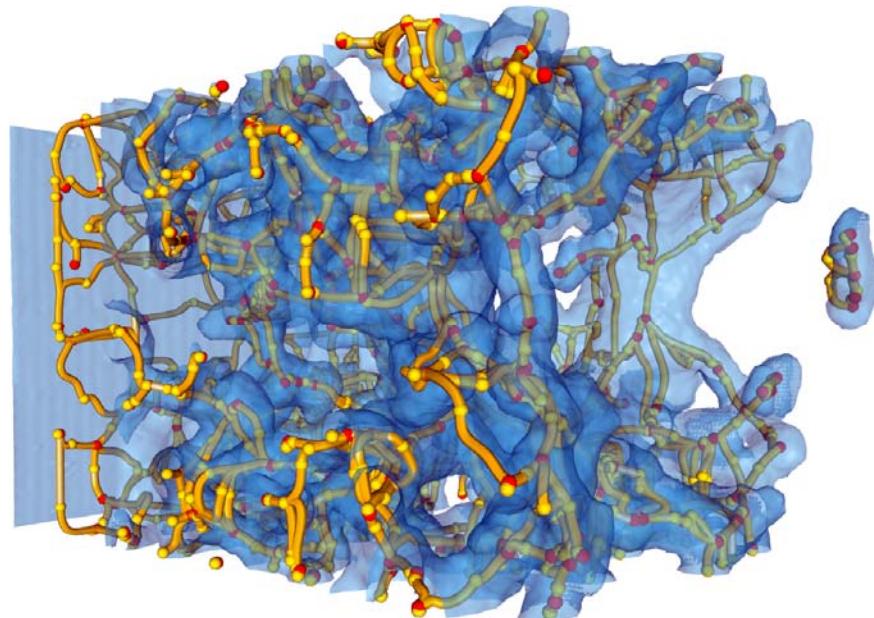
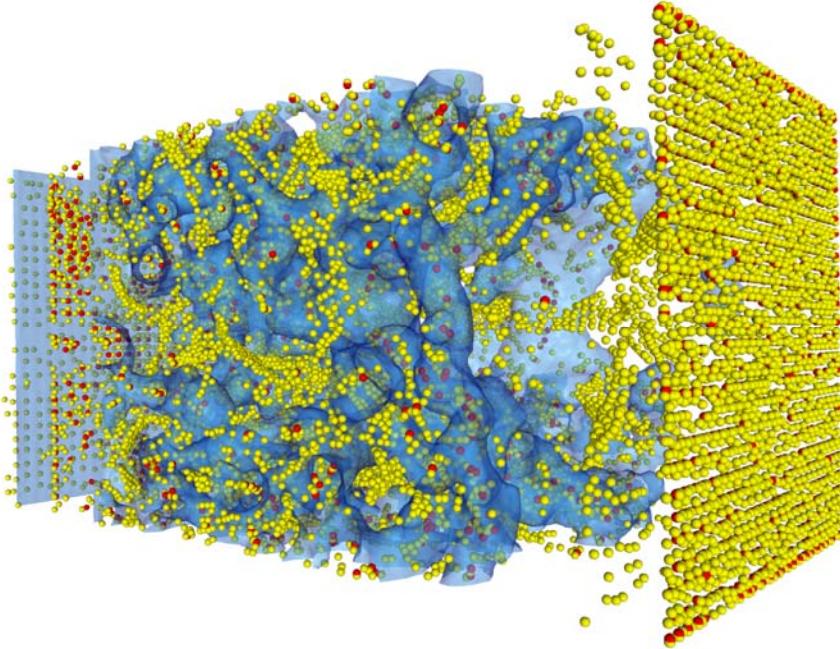


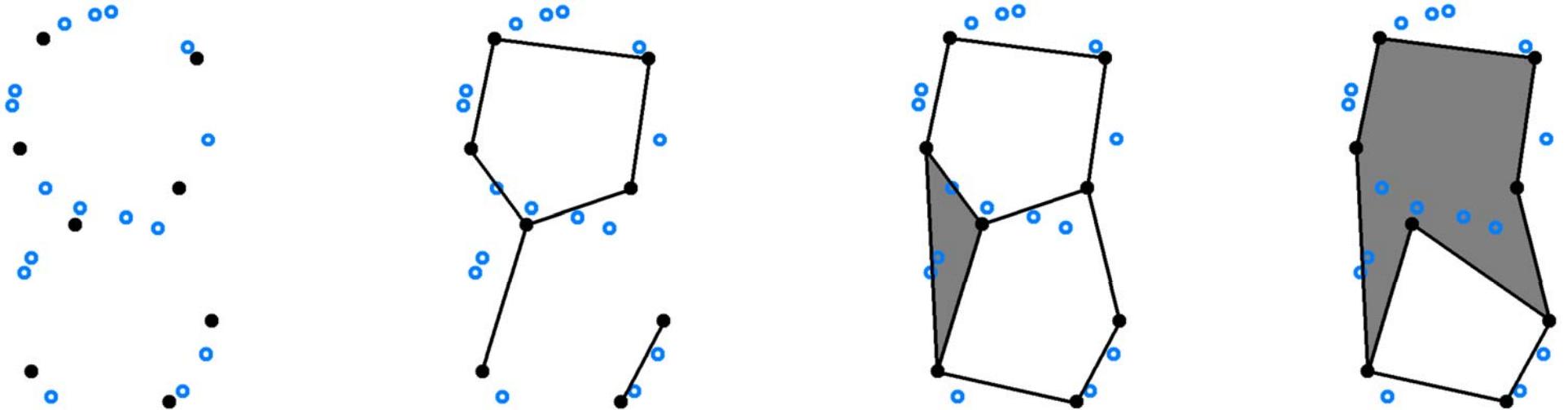
Sandia National Laboratories - CSRI Workshop on
Combinatorial Algebraic Topology (CAT): software, applications & algorithms
(by invitation only)
29–30 August 2009, Reception Evening 28 August
Hilton Santa Fe, NM

Interactive Data Analysis Session Using Morse-Smale Complexes

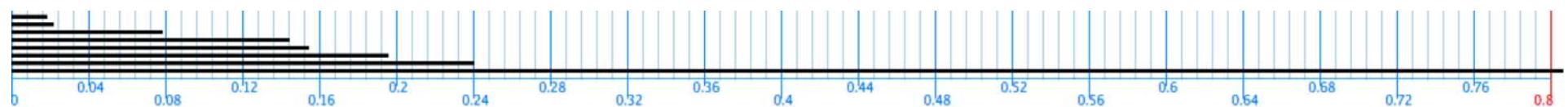


Multiple Topological Scales for Analysis of Porous Media

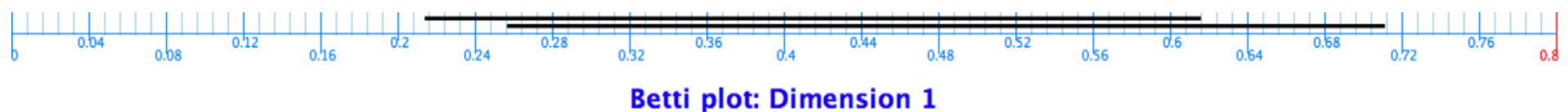




Betti barcodes for a figure-eight



Betti plot: Dimension 0

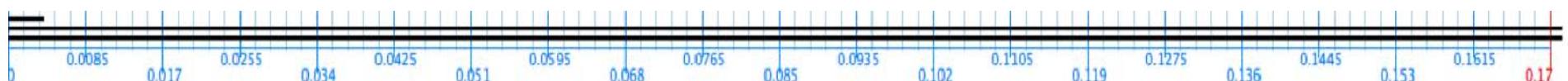


Betti plot: Dimension 1

Betti barcodes for a torus



witT100: Dimension 0



witT100: Dimension 1



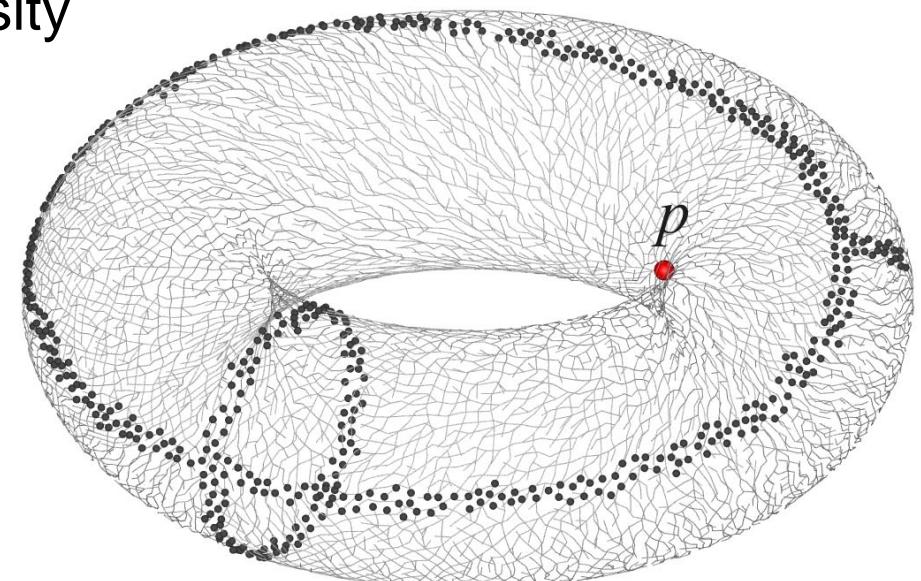
witT100: Dimension 2

Thomas Jefferson made barcodes too.

Cut Locus and Topology from Point Data

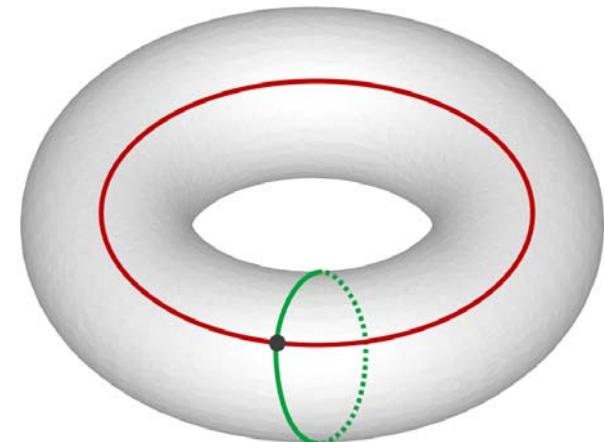
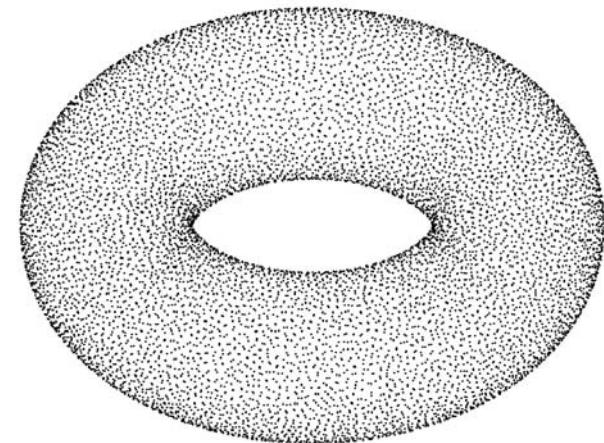


Tamal K. Dey Kuiyu Li
The Ohio State University



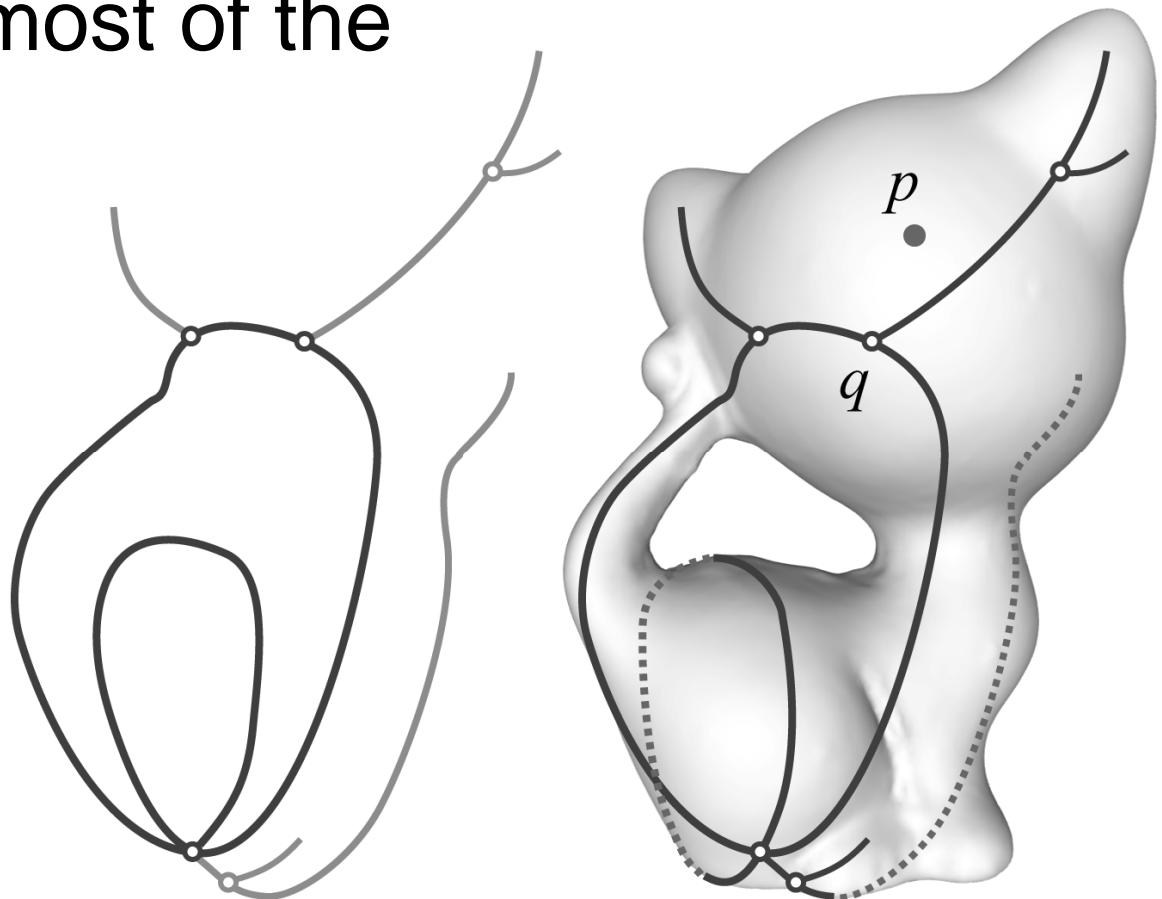
Problem

- Input: a point cloud P sampled from manifold M embedded in a higher dimensional Euclidean space \mathbb{R}^d
- Goal: extract topology of M from P



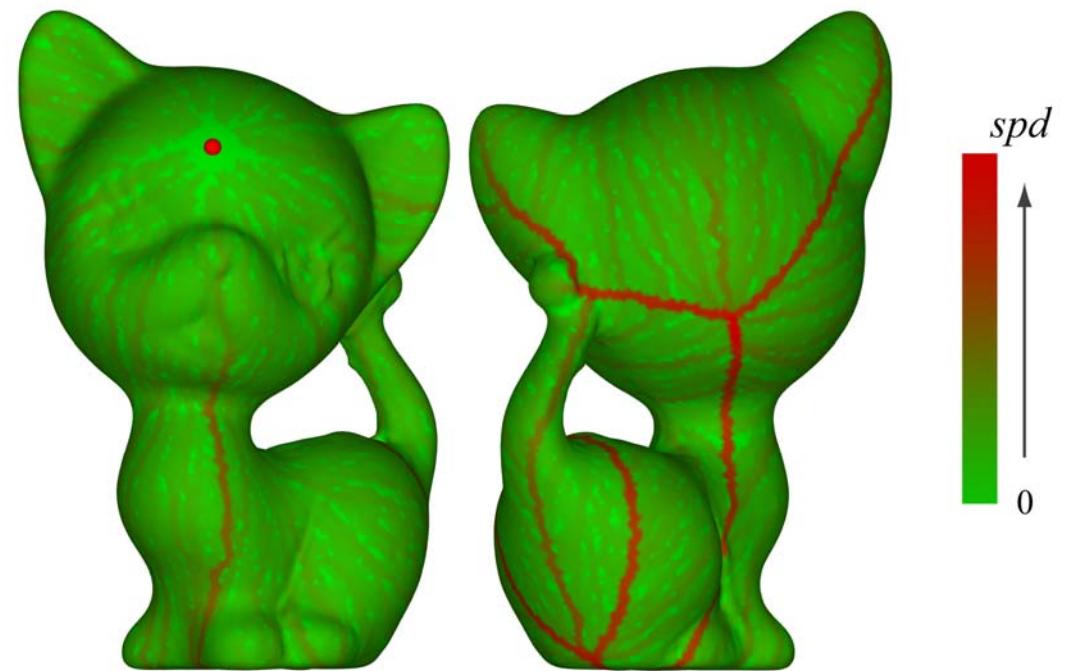
Cut Locus in M

- Observation: Cut locus of M retains most of the topology



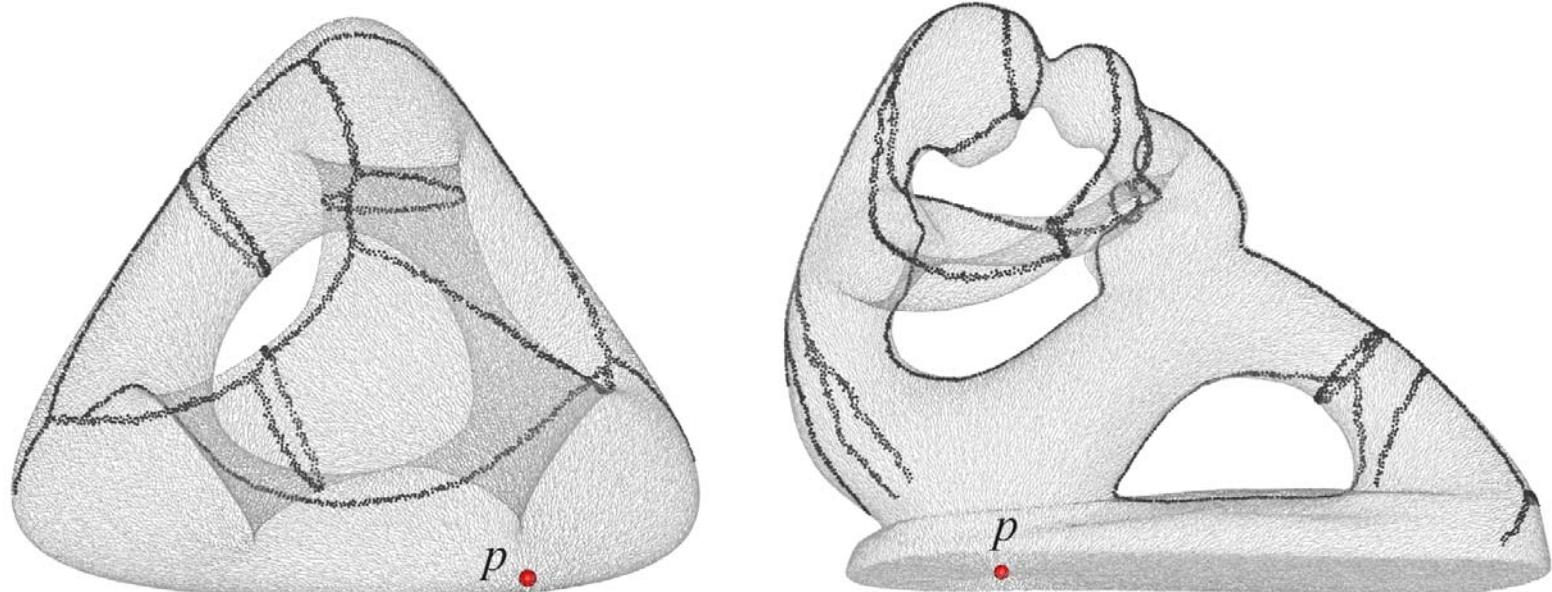
Cut Locus in P

- Use a small subset P' of P to approximate the cut locus of M
- P' is computed using the notion of **geodesic spread function**



Topology

- Use the method of [Chazal-Oudot 08] on P' to compute the Betti numbers of M
- Computationally efficient because P' is much smaller

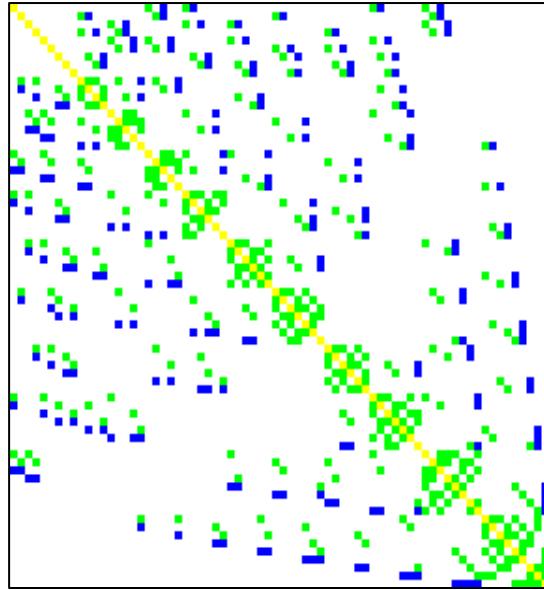


Paper

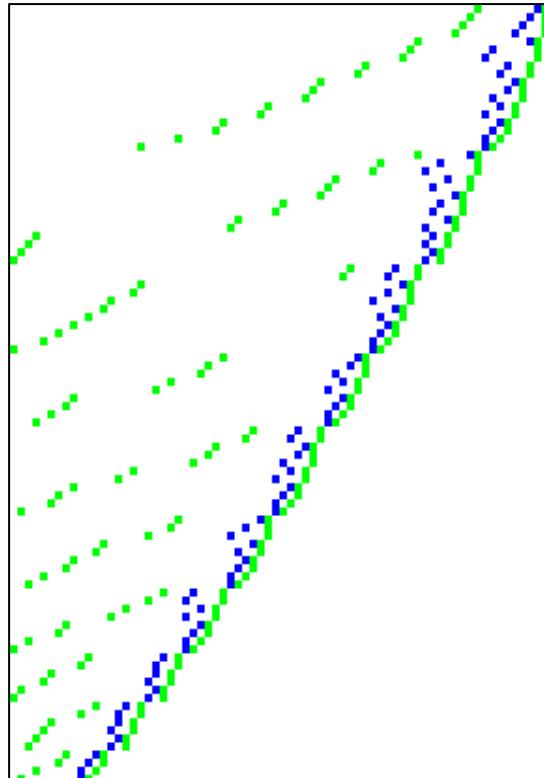
Cut locus and topology from surface point data. 2009. Tamal K. Dey, Kuiyu Li. Proc. 25th Ann. Sympos. Comput. Geom. (SoCG09), 125-134

Integer Smith form, ranks of Homology matrices

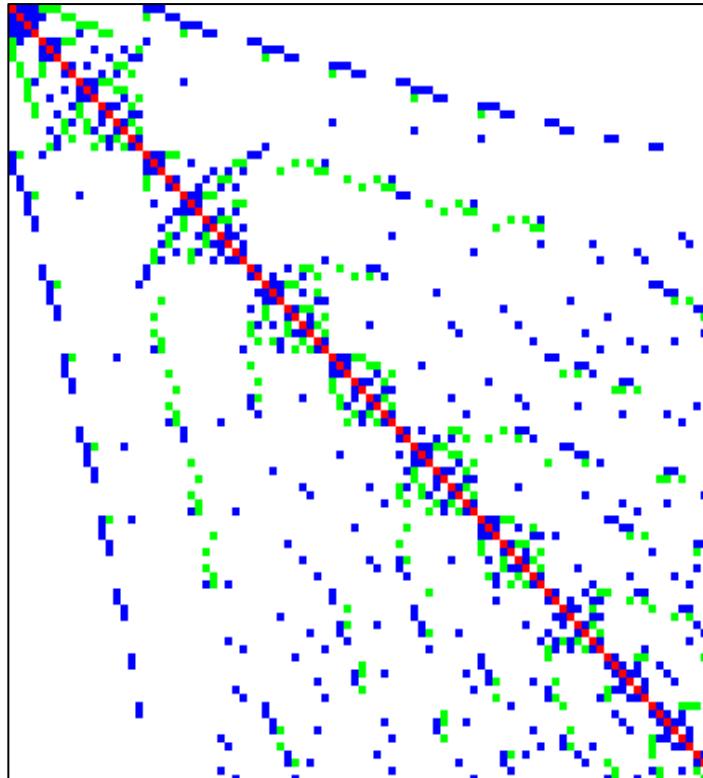
$A^T A$



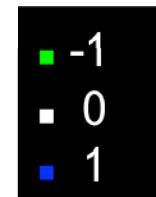
A



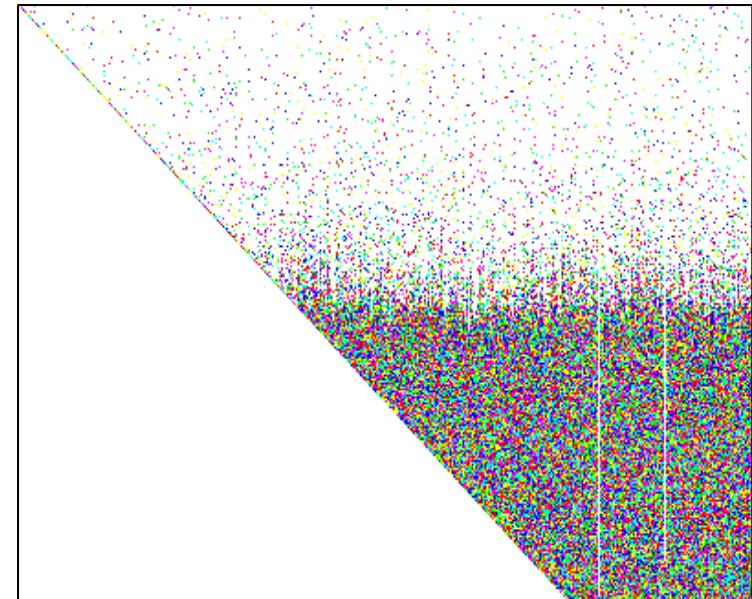
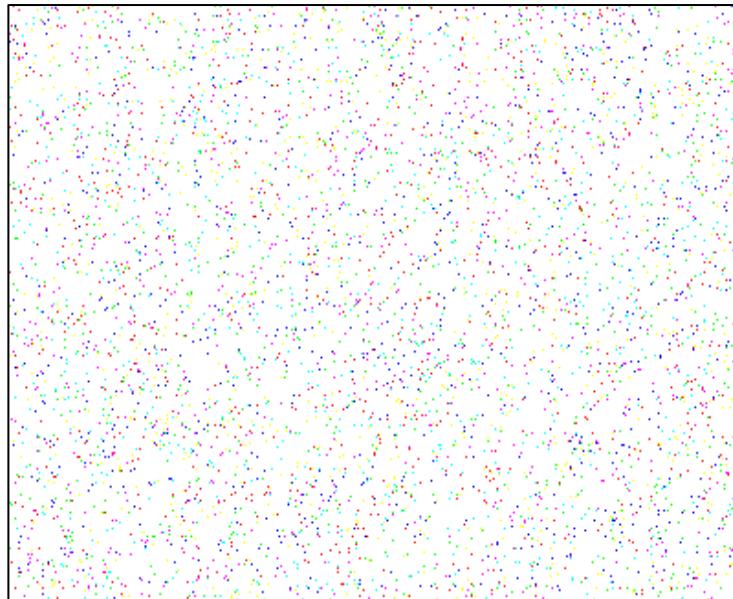
A^T



- Large dimensions
- Very sparse
- Small invariant factors
- Minimal polynomial of $A^T A$ has very small degree



Direct or iterative methods ?



Fast at start, Fat at last

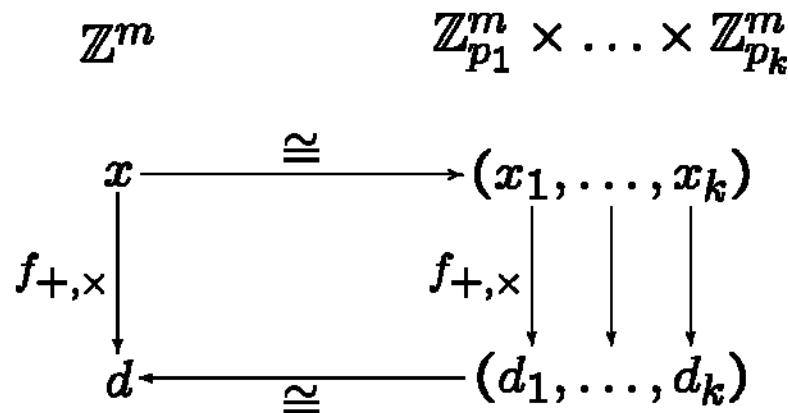
$$y \in F^n \rightarrow \boxed{A} \rightarrow Ay \in F^m$$

Slowly but Surely

Integer coefficient growth

$A := [[2, 1, -3, 7, -4, 0, 9, 1, 6, -8, [6, -2, -8, -7, 4, -5, 5, 2, 4, -9, [-1, -7, 2, 9, 2, 3, -1, 6, -8, 7, [3, -2, 0, -4, 8, 1, 2, 4, 6, 6, [-4, -6, -3, -7, -8, 2, 4, -1, -2, -9, [8, 7, 6, -3, -9, 3, 4, 6, 3, 6, [-6, -8, 3, -3, -1, -5, 6, 2, 2, -5, [-5, -2, -9, -6, -1, -3, 5, 8, 8, -1, [2, 3, 8, -9, 3, -1, 8, -6, 7, -7, [-8, 3, 7, 4, 3, 2, 2, -9, -2, -4]]];$, and $b := [-2, -9, -3, 8, 0, -2, -6, 8, 5, -7]$;
→ Solution of $Ax=b$:
 $x = [-154378467, 166174189, 556790464, 72535956, 258019417, 499902615, -930948611, 1007780694, 356546217, -1001276698] / 55181746$

Chinese remaindering, Hensel lifting

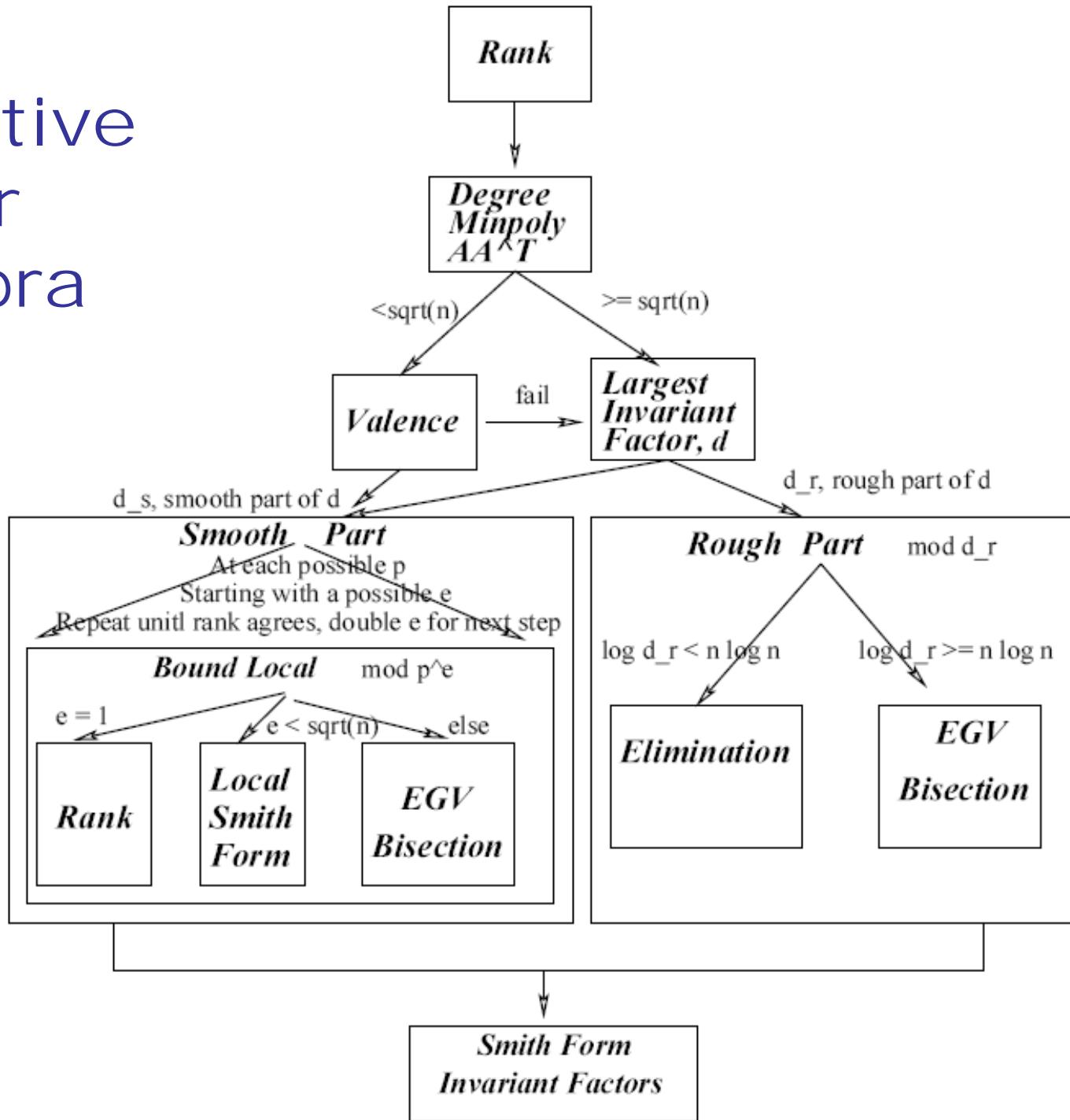


$O(Pb \cdot k + mk^2)$, instead of $O(Pb \cdot k^2)$

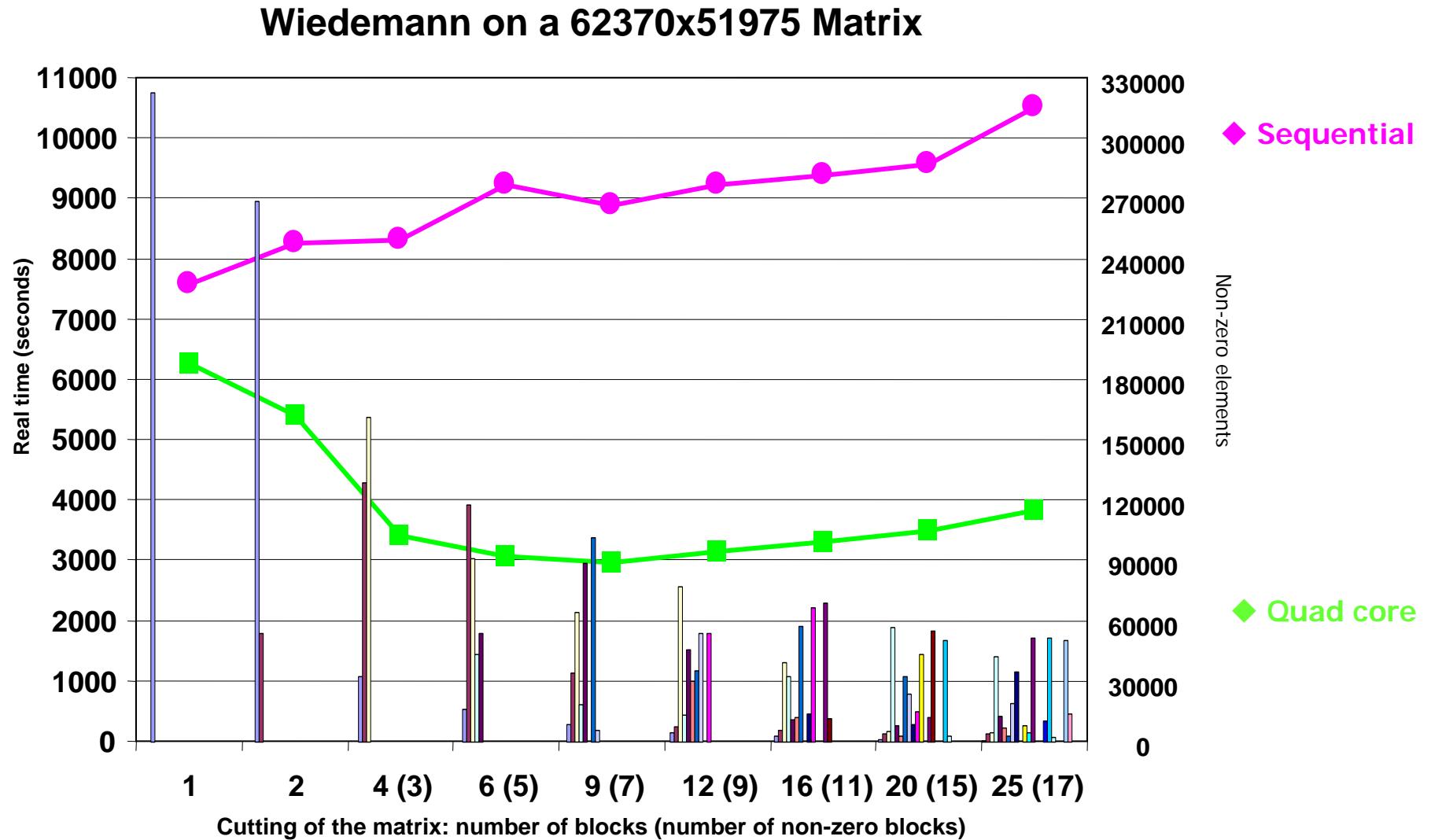
- | | |
|--|----------------|
| 1. PLUQ(A) mod p | n^w |
| 2. Iterate k times ($k < n/2 \log(n)$) | |
| • 2 triangular solve $Ax_i \cdot b_i \bmod p$ | $2n^2$ |
| • $b_{i+1} = (b_i - Ax_i)/p \bmod Z$ | $4n^2 \log(n)$ |
| 3. $x = x_0 + p x_1 + p^2 x_2 + \dots + p^k x_k$ | k^2 |
| 4. Rational reconstruction | nk^2 |

$O(n^w + n^3 \log^2(n))$, instead of $O(n^{w+1} \log(n))$

Adaptive linear algebra

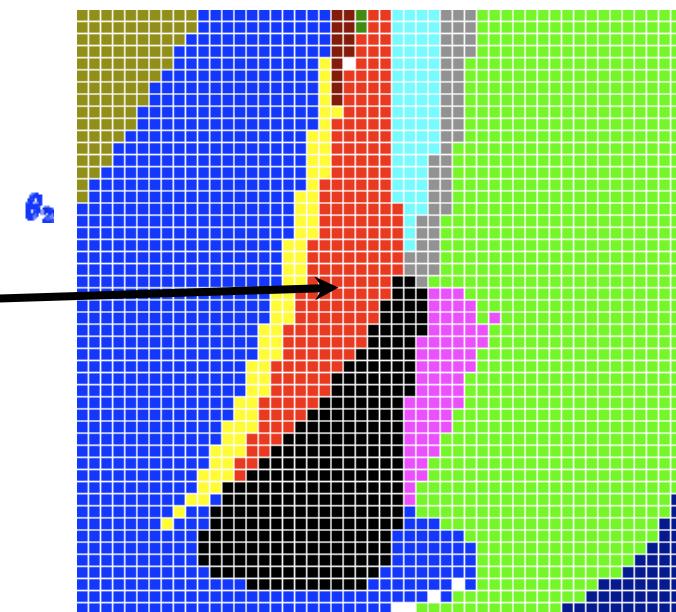
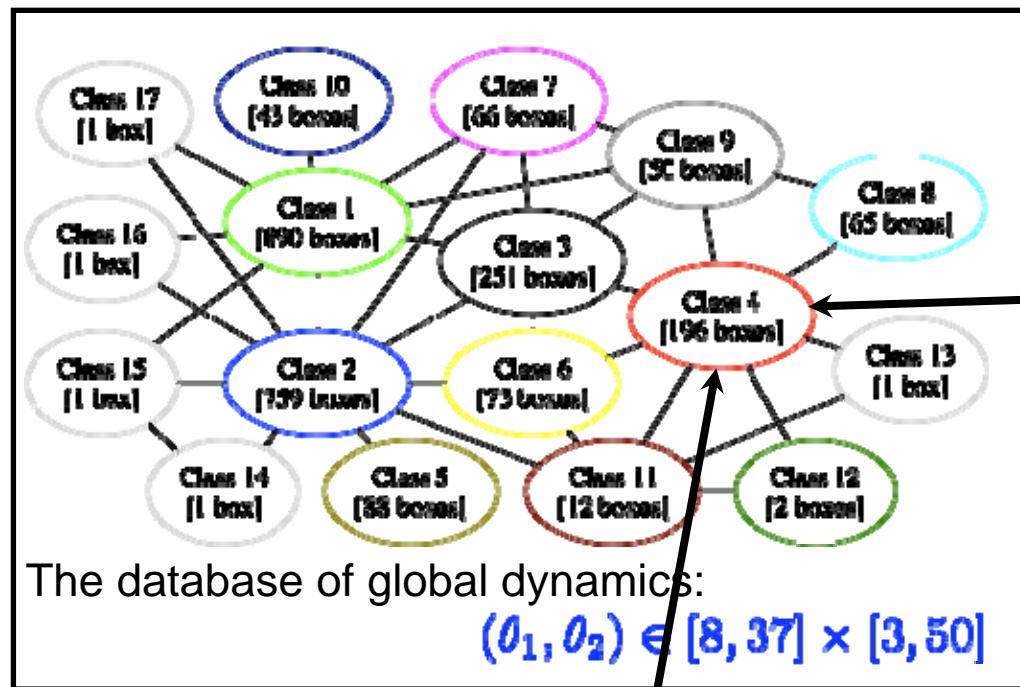


Scalability: multicore algorithms



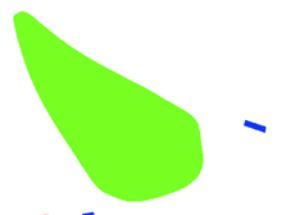
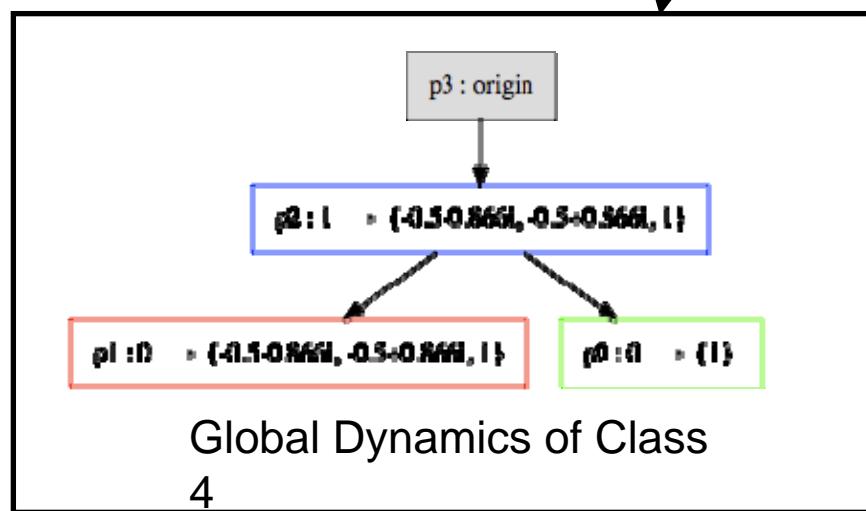
A Two Parameter Population Model:

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \mapsto \begin{bmatrix} (\theta_1 x_1 + \theta_2 x_2) e^{-0.1(x_1+x_2)} \\ 0.7x_1 \end{bmatrix}$$



Continuation Classes in Parameter Space

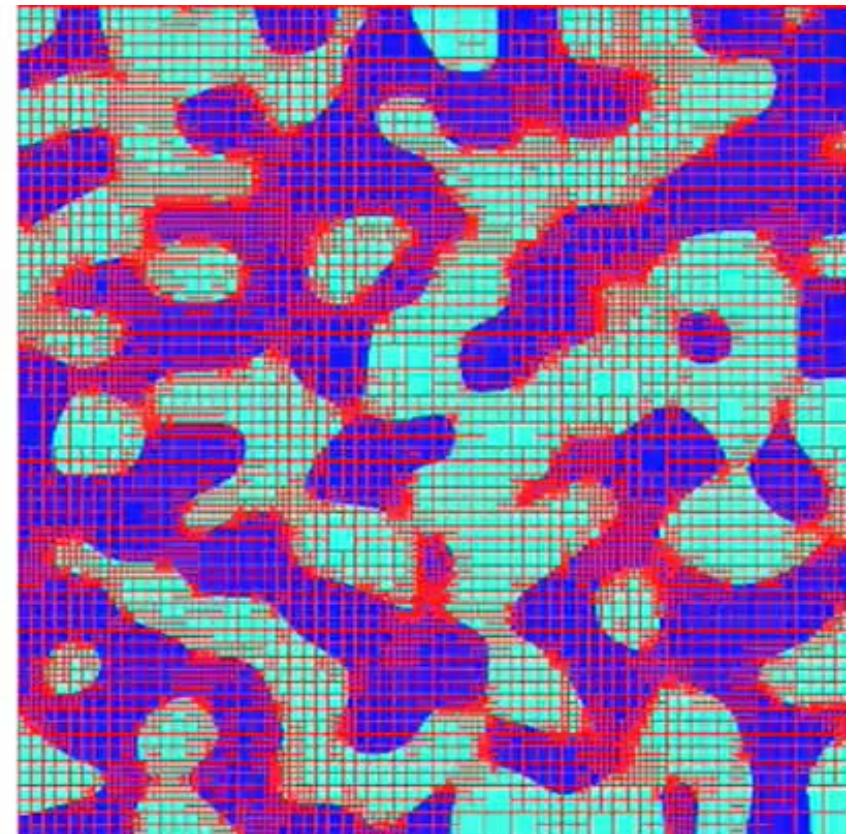
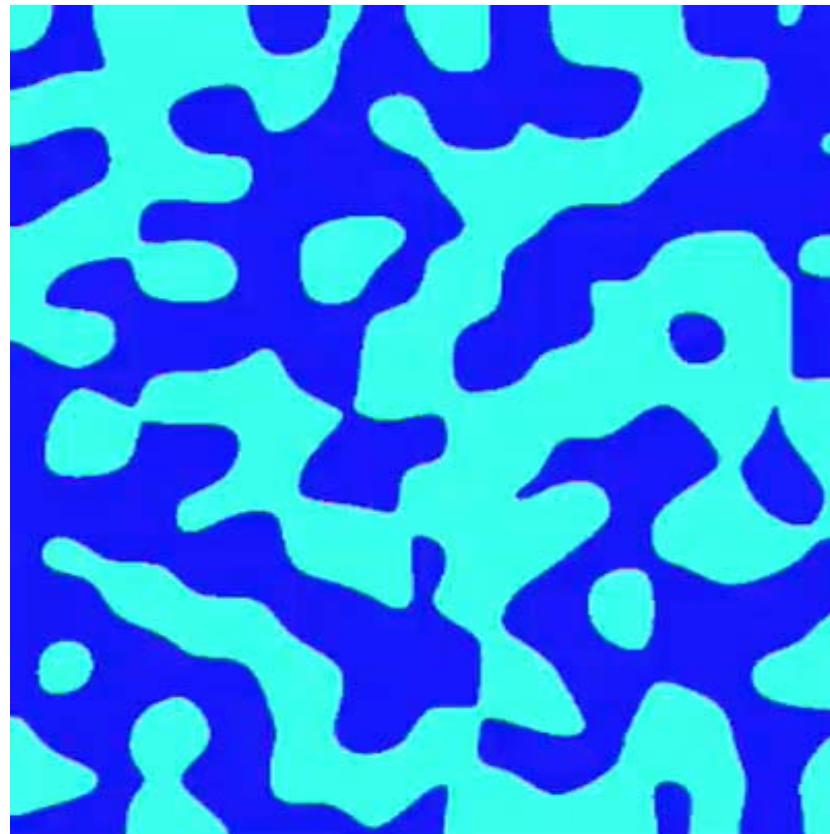
A Database Schema for the Analysis of Global Dynamics of Multiparameter Systems,
 Z. Arai, W. Kalies, H. Kokubu, K. Mischaikow,
 H. Oka, P. Pilarczyk, SIADS, 2009



Dynamics at a particular parameter value in red region.

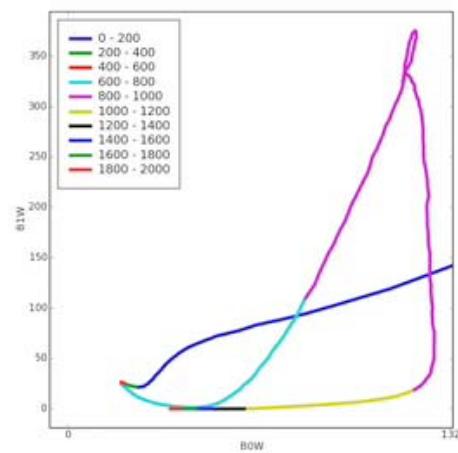
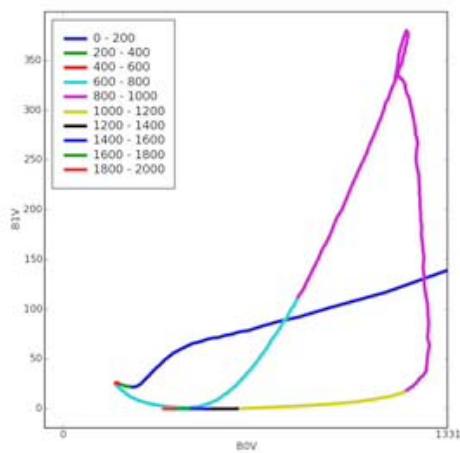
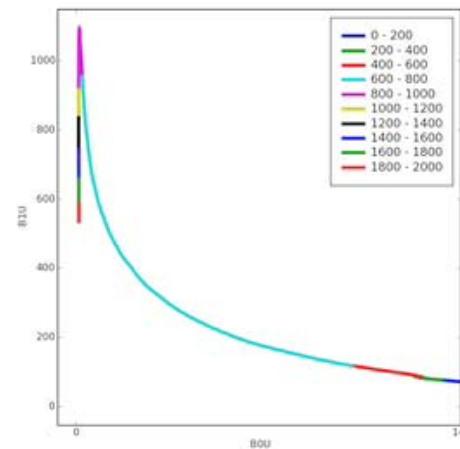
An adaptive cubical verification algorithm for the homology of evolving nodal domains

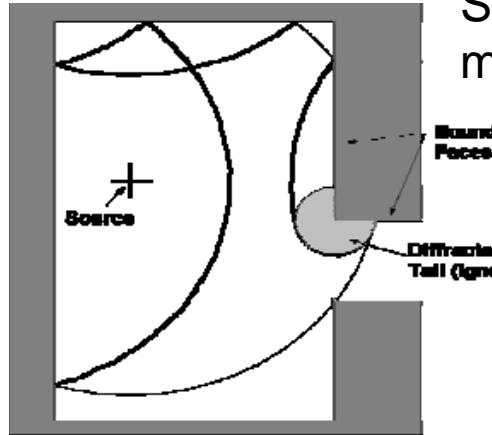
Day, Kalies, Wanner (2009)



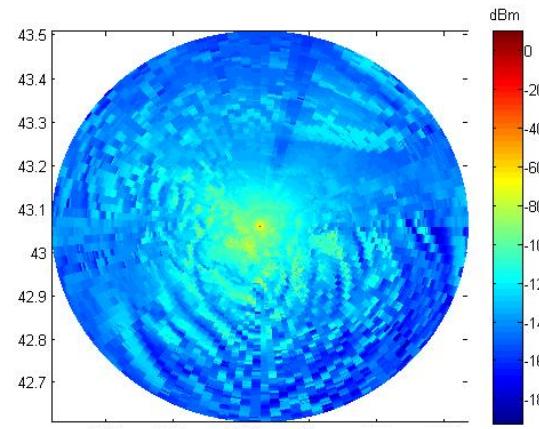
Pattern morphology evolution in ternary alloys

Corrigan, Mischaikow, Wanner (2009)



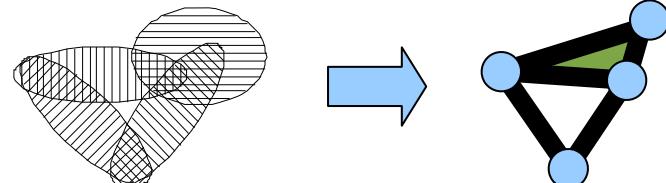
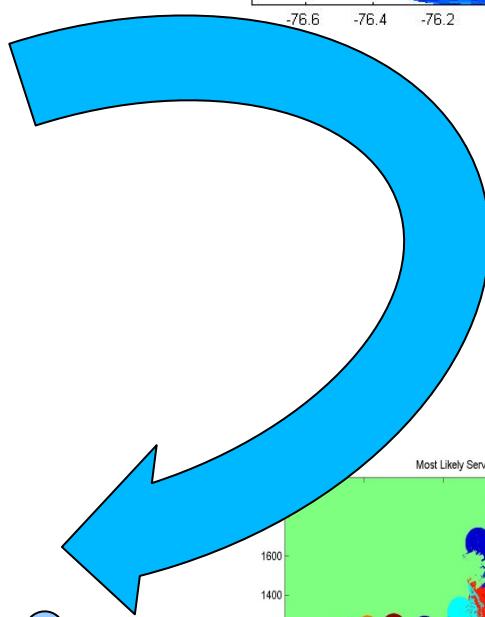


Simplified wave models

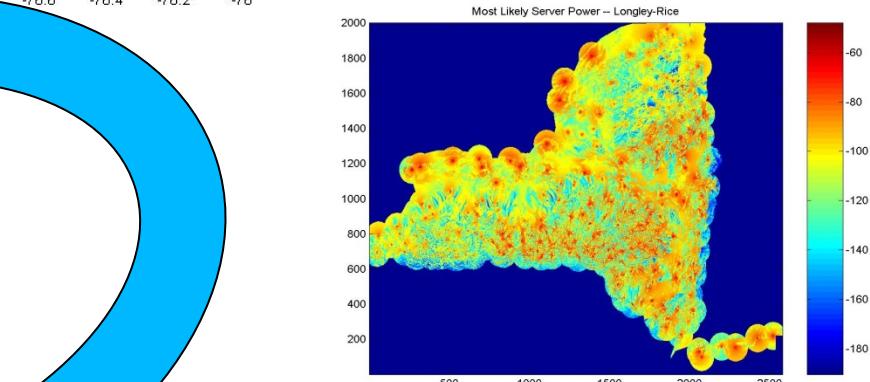


Propagation modeling for transmitters

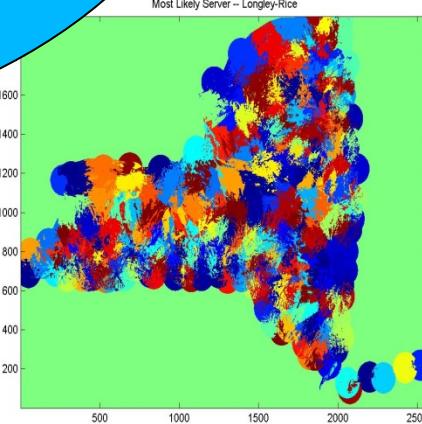
Topology for communications network performance and capacity analysis



Interference complexes and sheaf-based approaches for network analysis

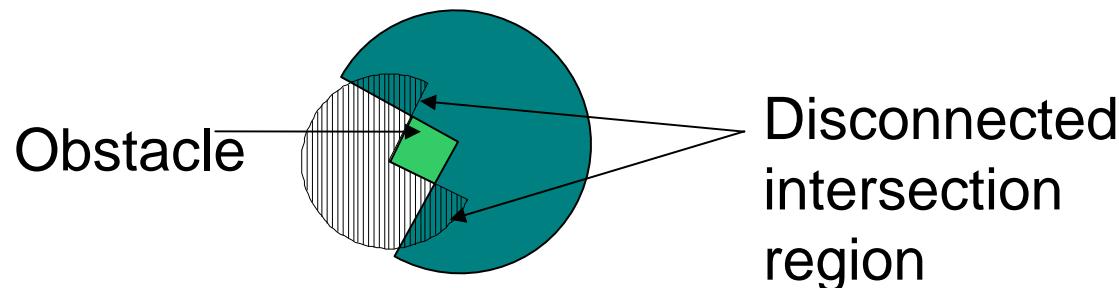


Full network and interference modeling

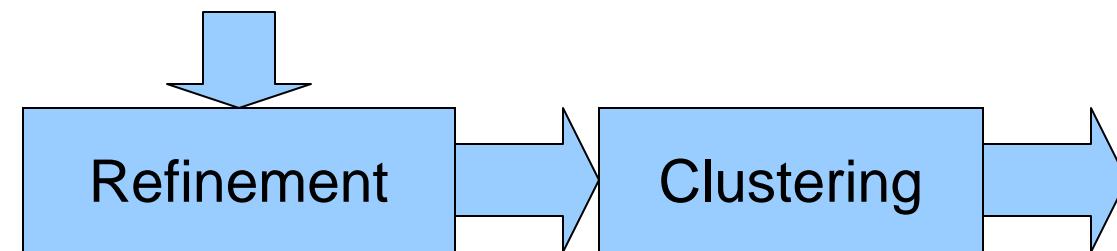


M. Robinson
 Penn

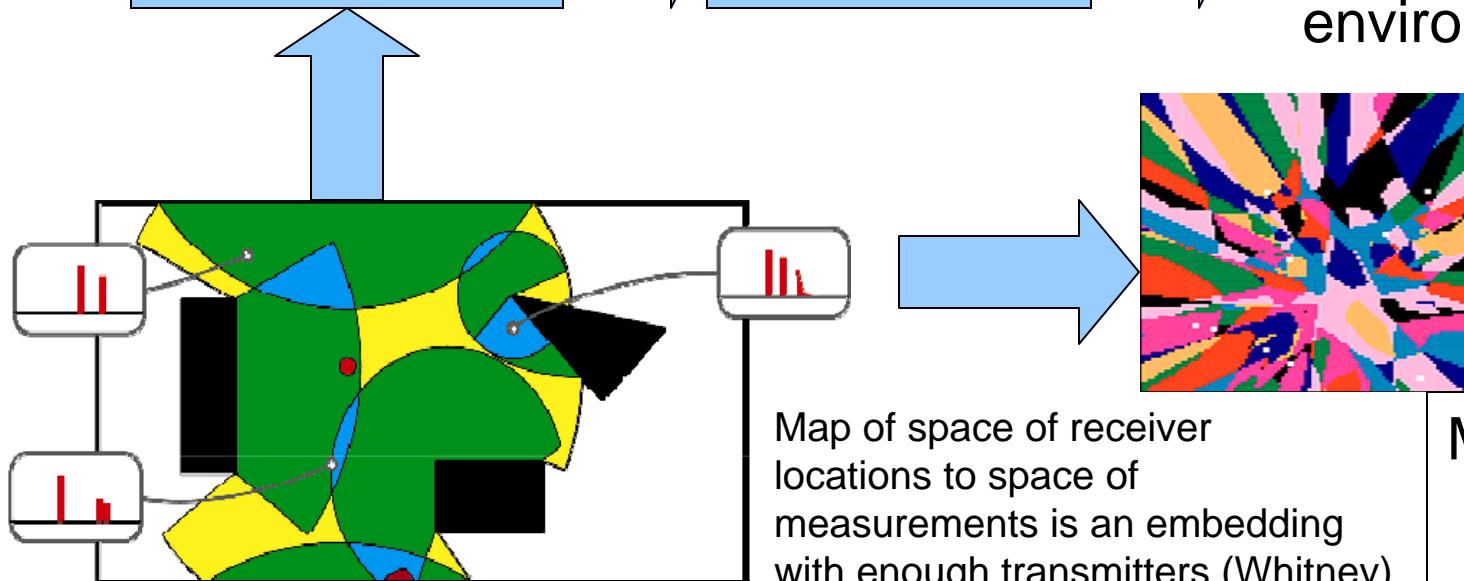
Discrete Opportunistic Topological Imaging



Regions of impact for opportunistic transmitters forms a cover



Recovers homotopy type, geometric information of propagation environment



Localization of receivers

M. Robinson
 Penn

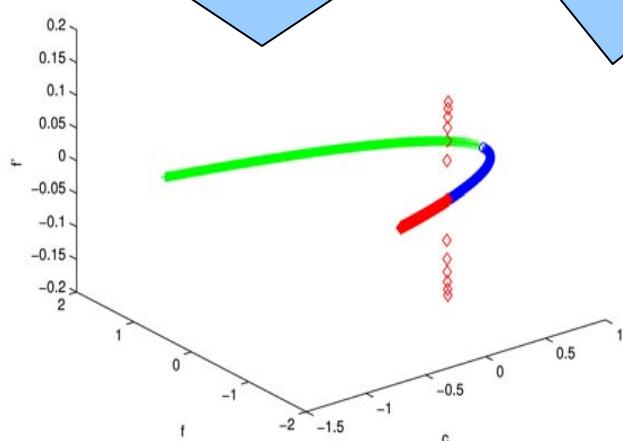
Semilinear parabolic
equation

Dynamical properties of
global solutions

Equilibrium analysis

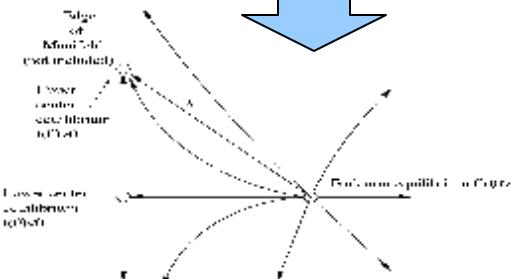
Evolution of stability

A Floer
theory for
semilinear
parabolic
equations



Bifurcation diagram

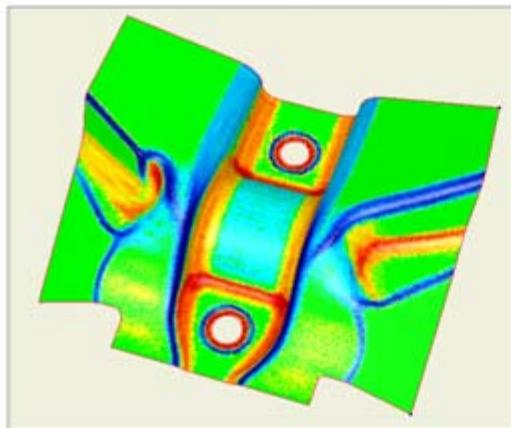
Cellular complex
of global solutions



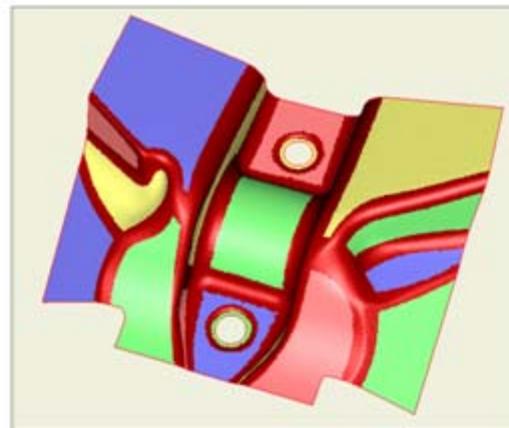
A topological index

M. Robinson
 Penn

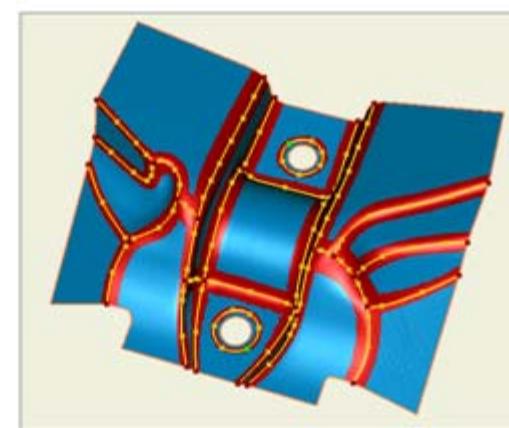
Steps in Functional Decomposition / 1



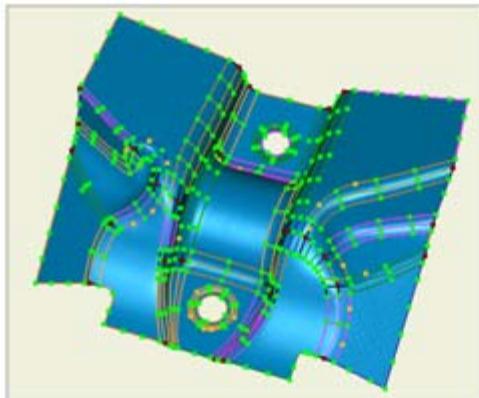
a. Numerical Curvature Map



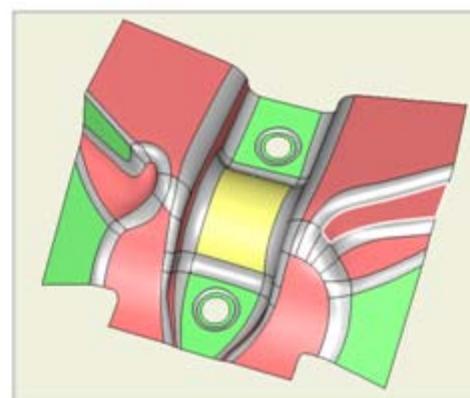
b. Regions and Separator Sets



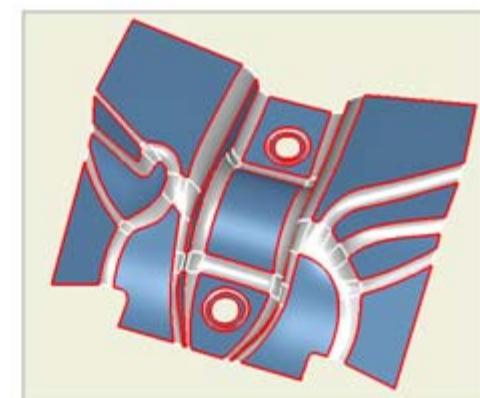
c. Contour extraction



d. Extending Contours

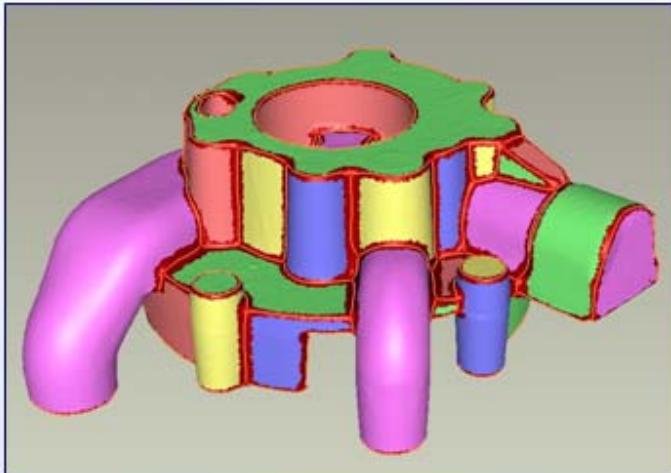


e. Classification

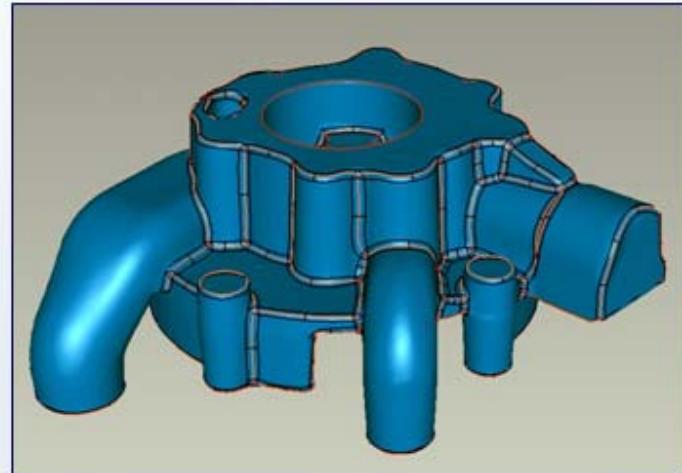


f. Trimmed Primaries

6. Automatic (rapid) surface generation II.



Segmented regions, separator sets
and feature skeleton (contours)



Thickened feature skeleton and
connecting features

Segmentation using Morse theory

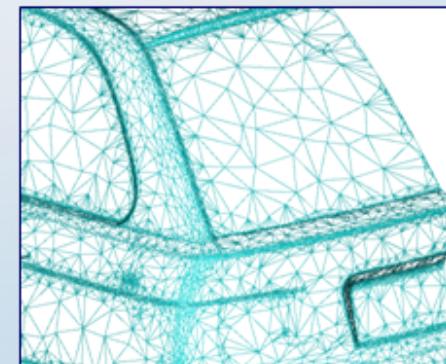
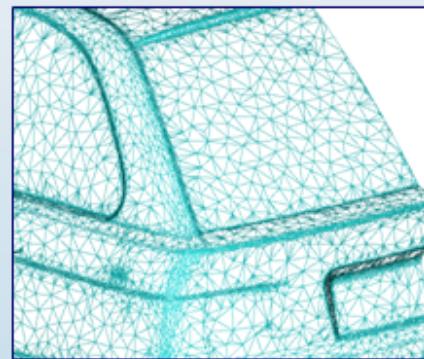
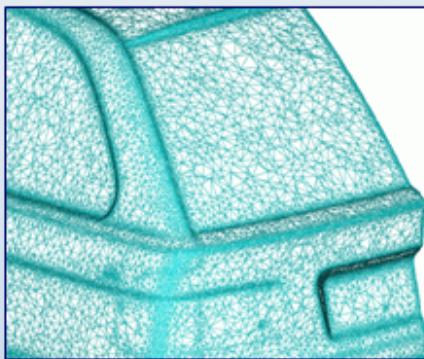
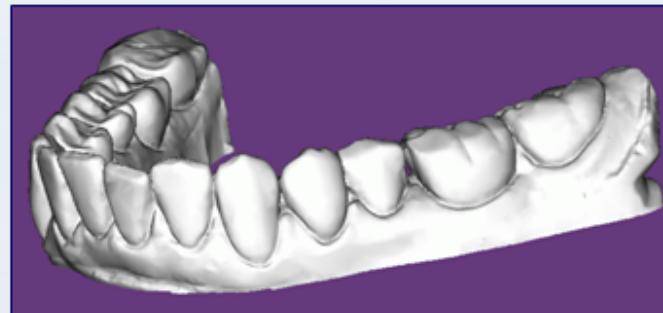
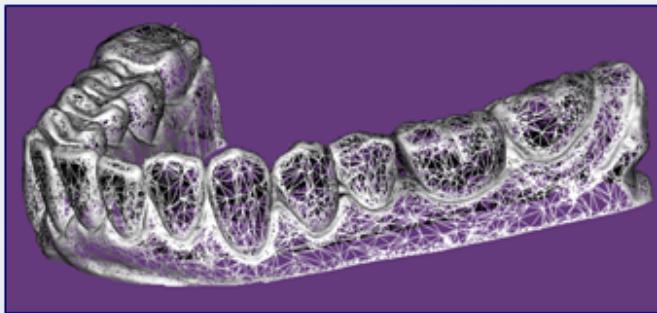
- smooth function over a manifold
- critical points - minimum, saddle, maximum
- piecewise linear Morse complex
- topological persistence and simplification

[Edelsbrunner et al 02](#), [Facello et al. 04](#), [Edelsbrunner 05](#)

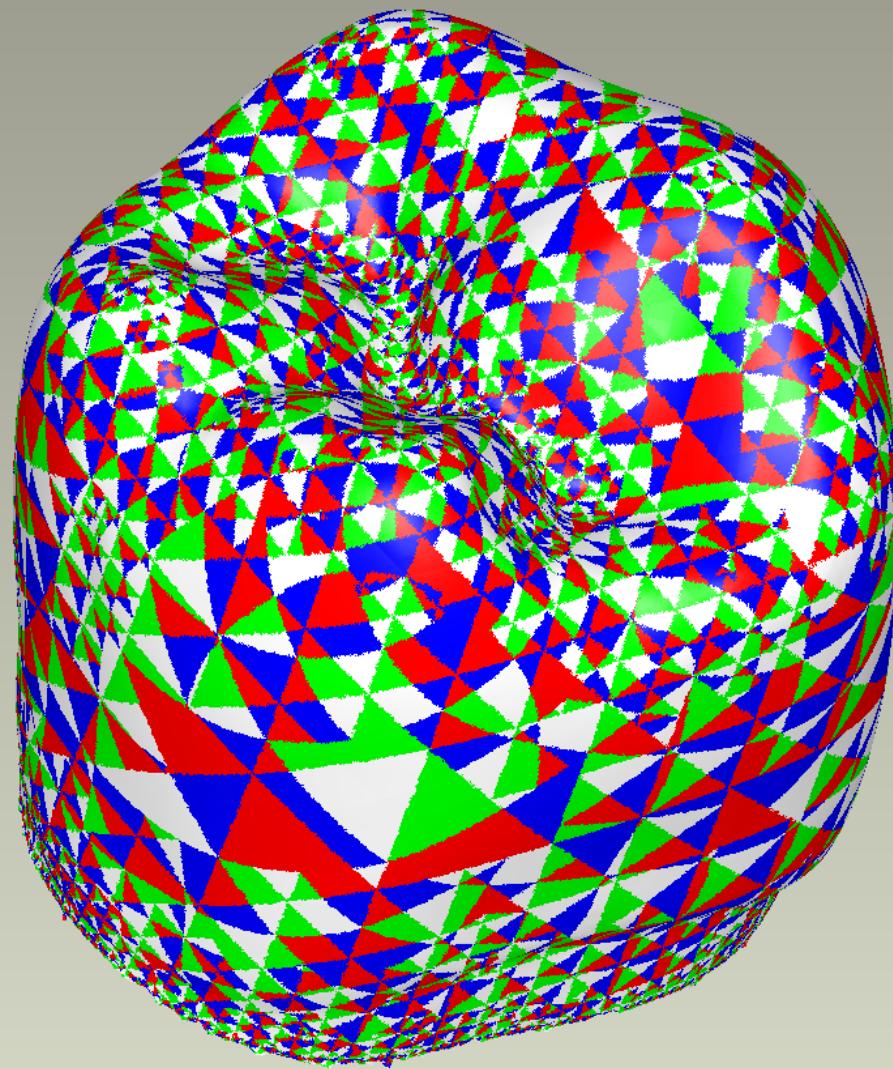
Triangulation/Decimation

well-established techniques

Hoppe et al. 92; Edelsbrunner, Mücke 94; Kos 01; Owen 98;
Taubin et al. 98; ...

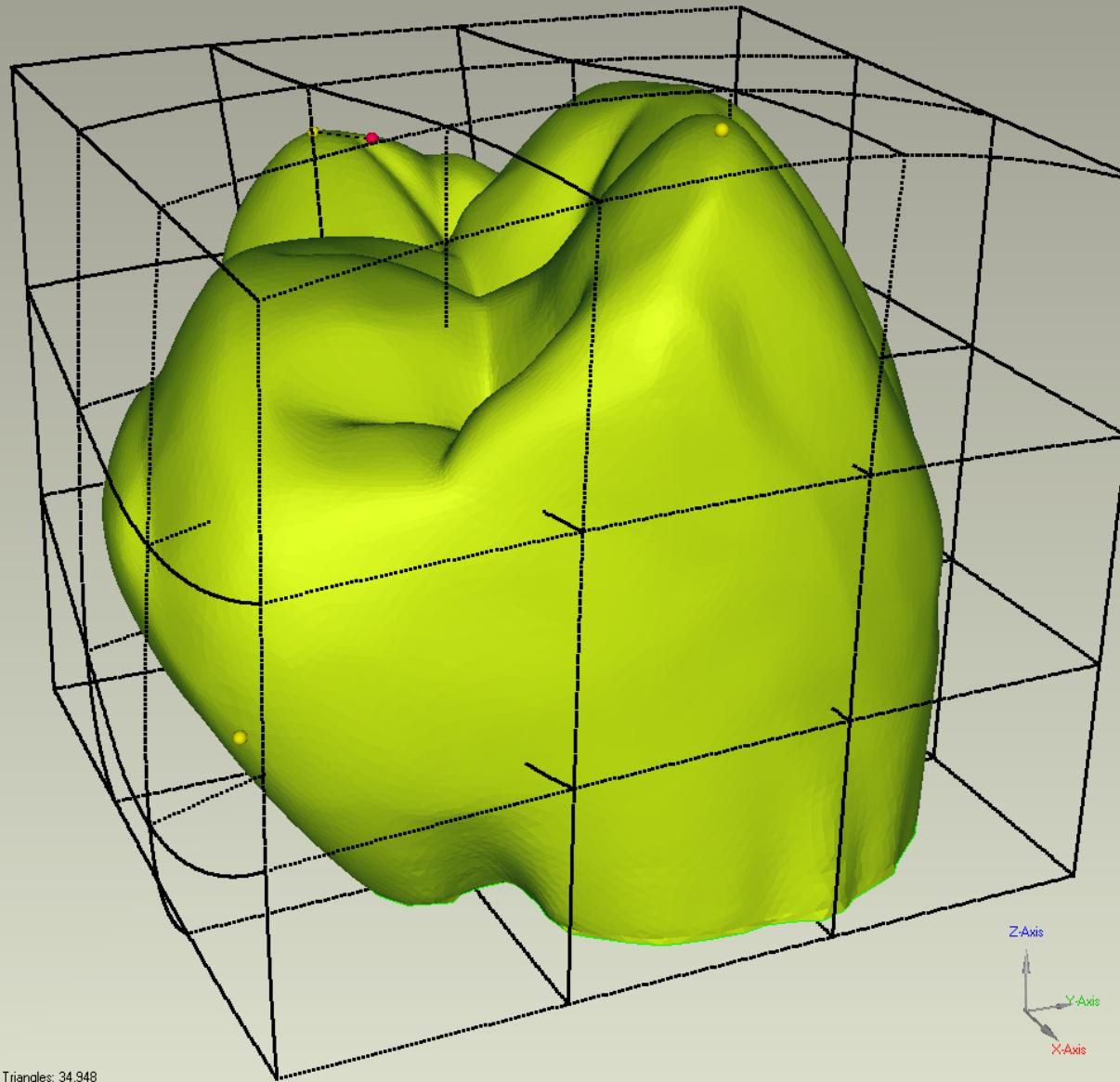


[Left]

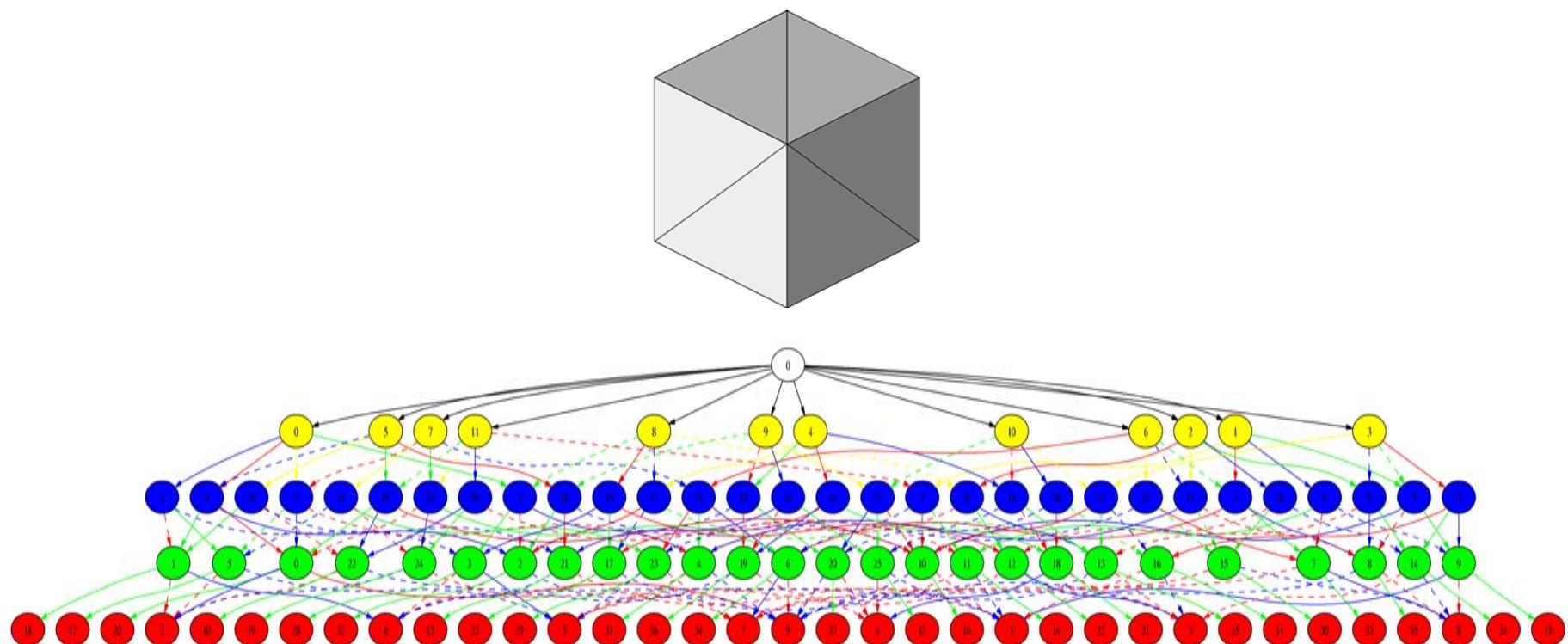


Current Triangles: 1,470,069
Selected Triangles: 0

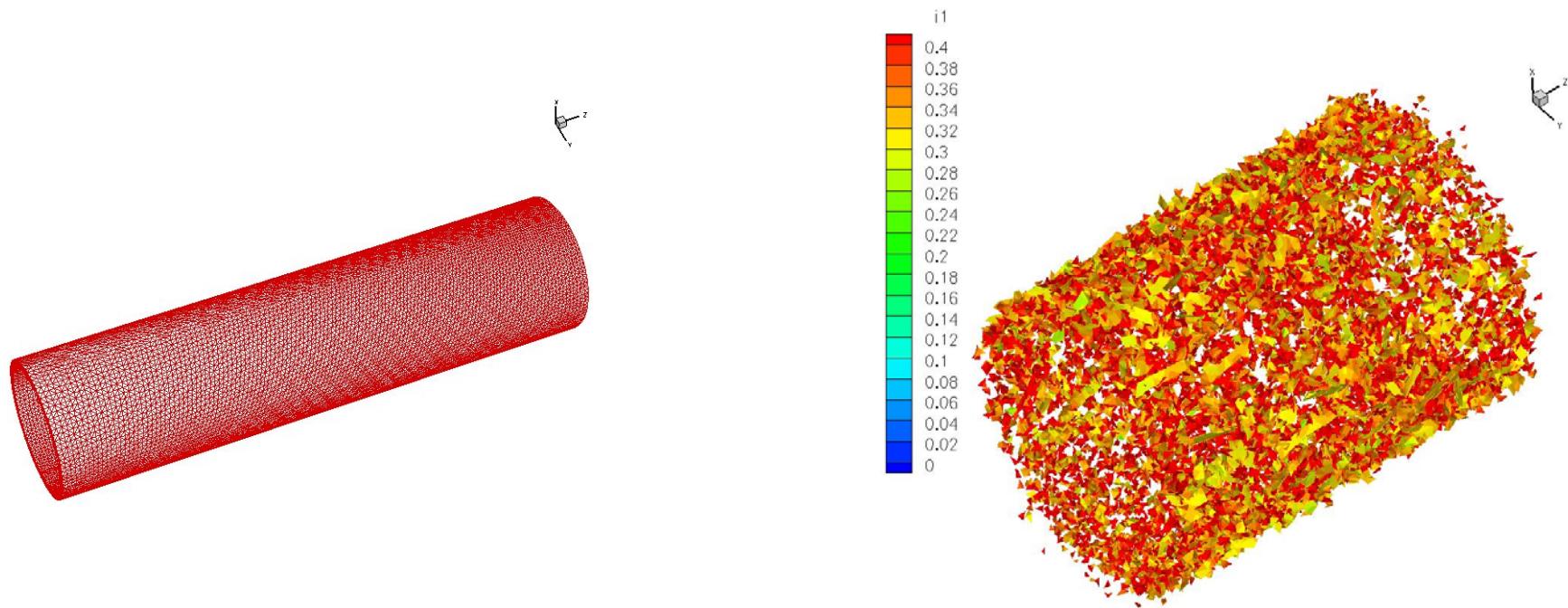
(Isometric)



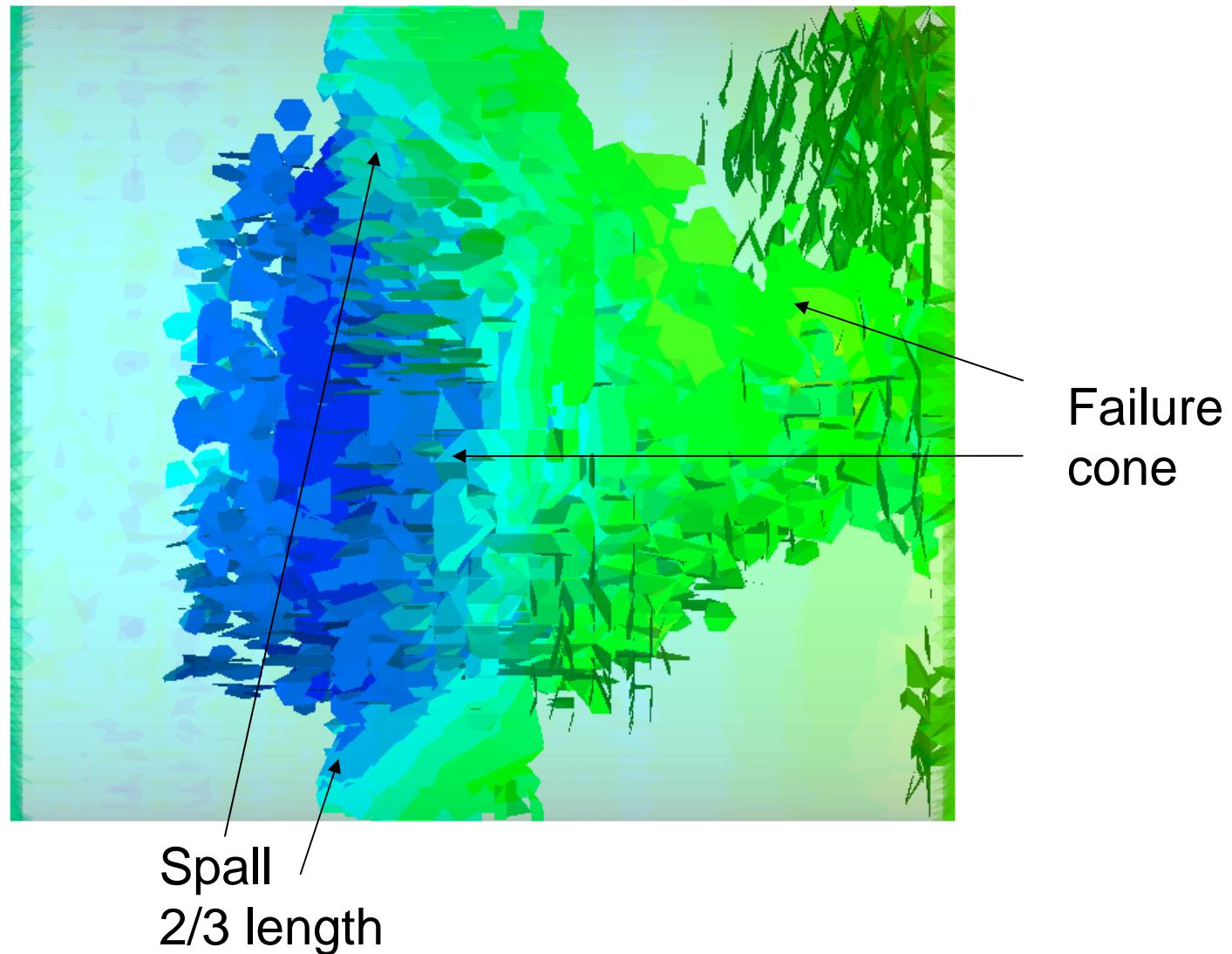
Mesh of Cube as Graph

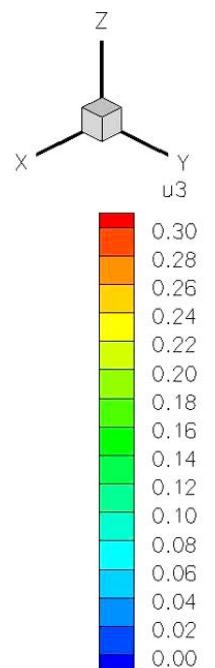


Testing of Graph Fracture

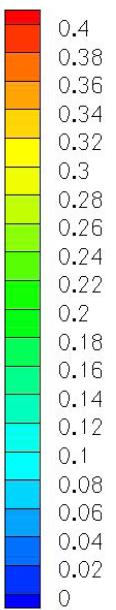


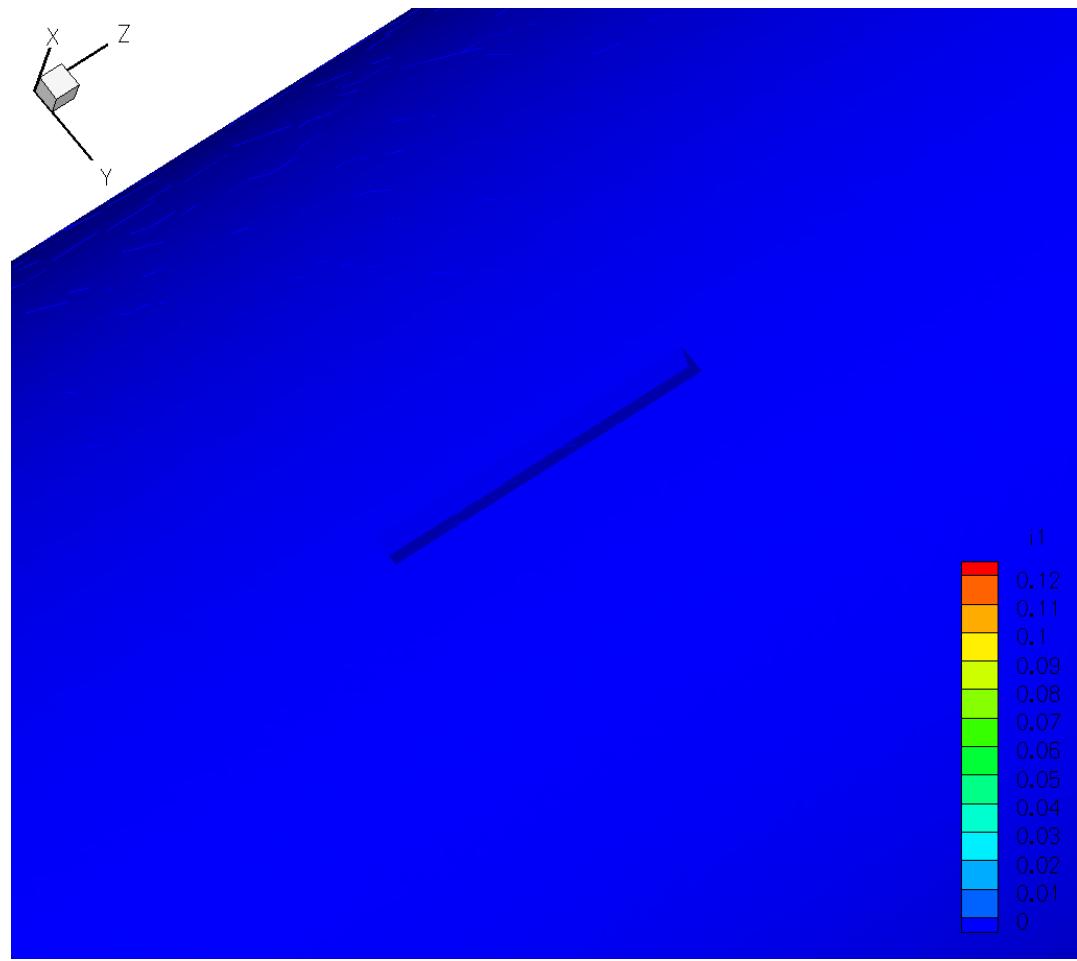
Fracture in Gypsum Cylinder

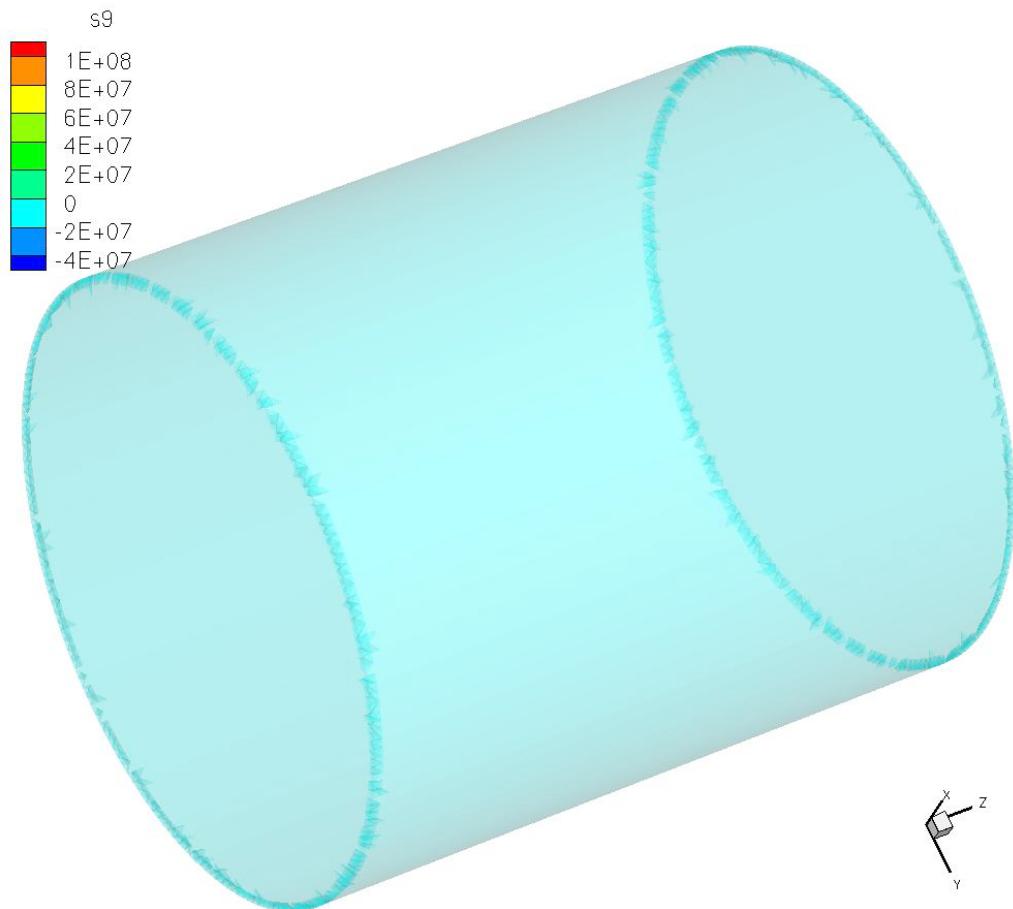




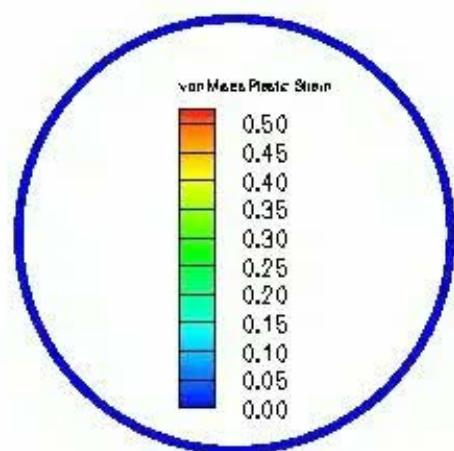
i1

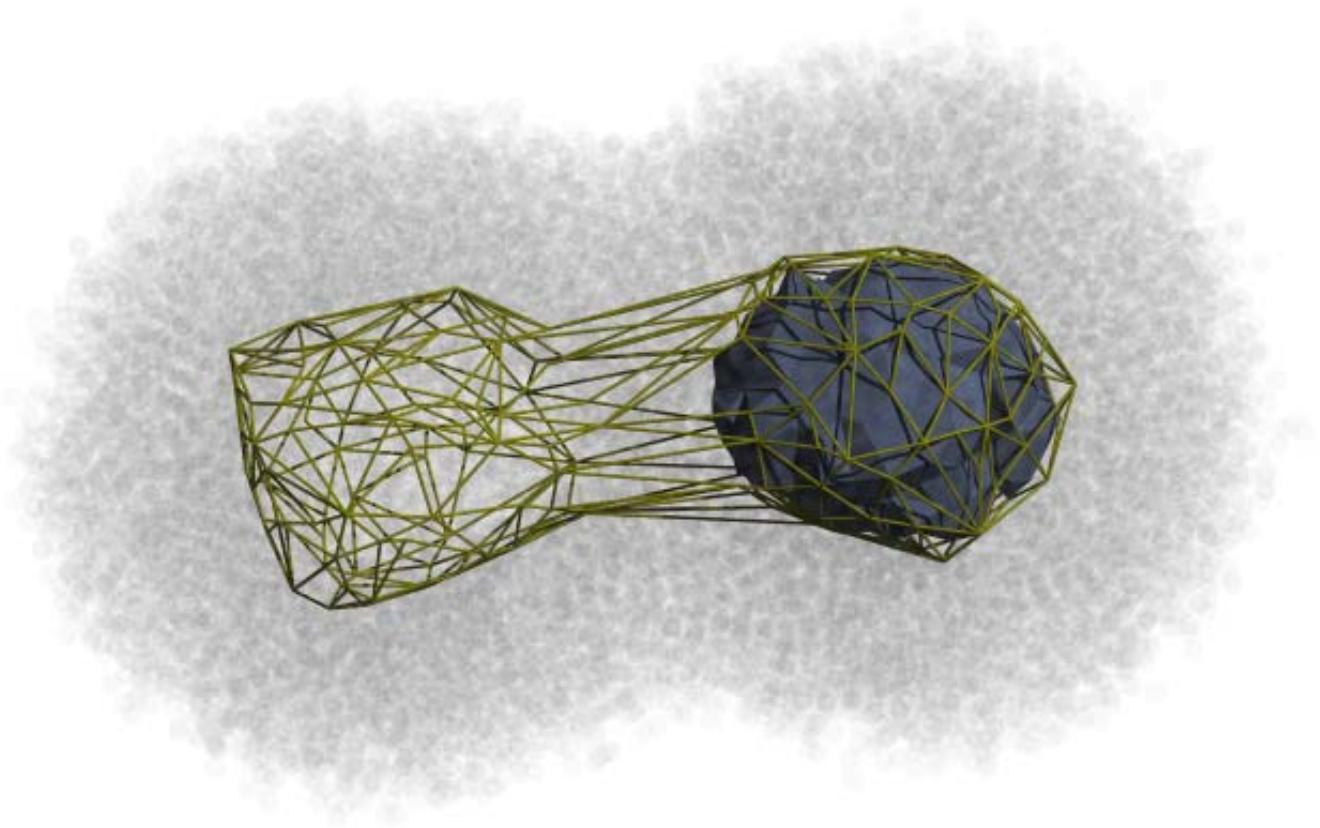


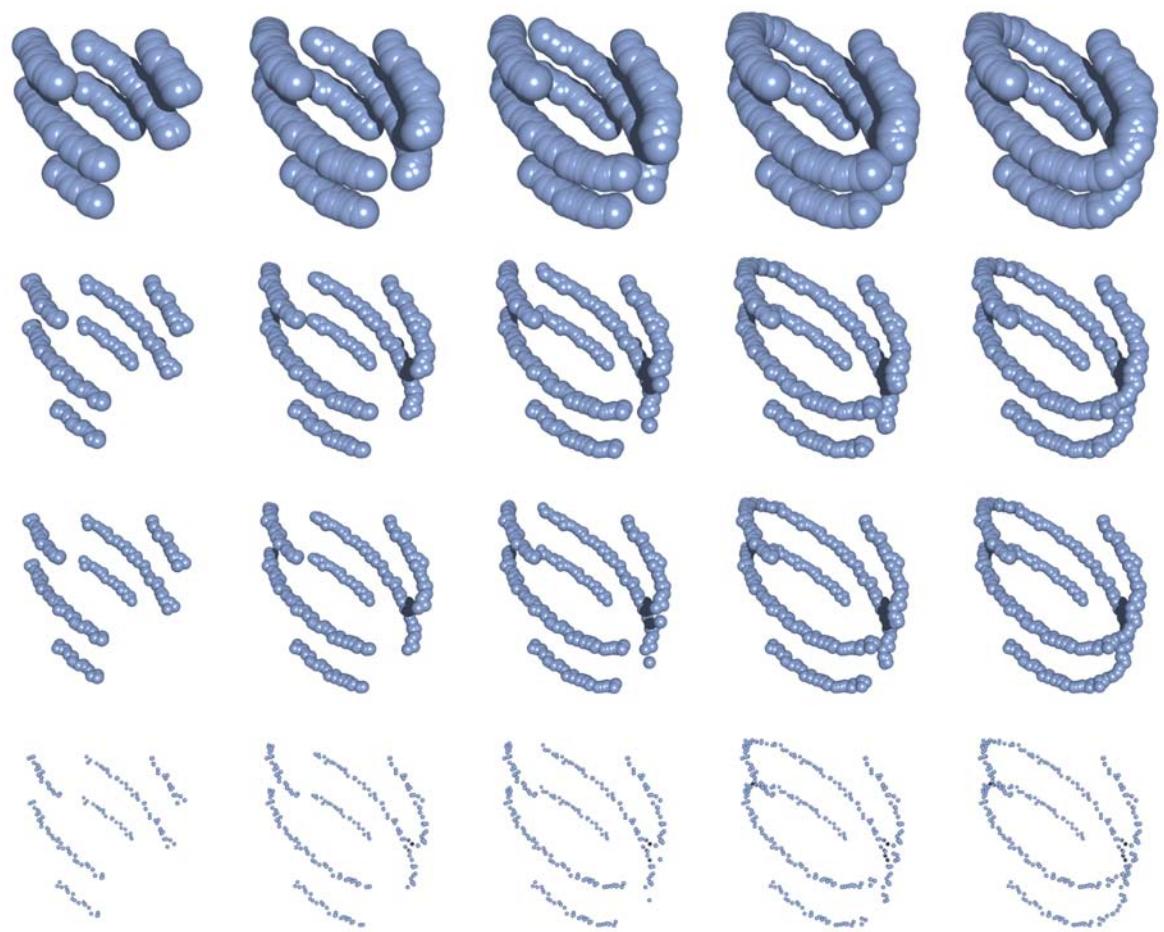


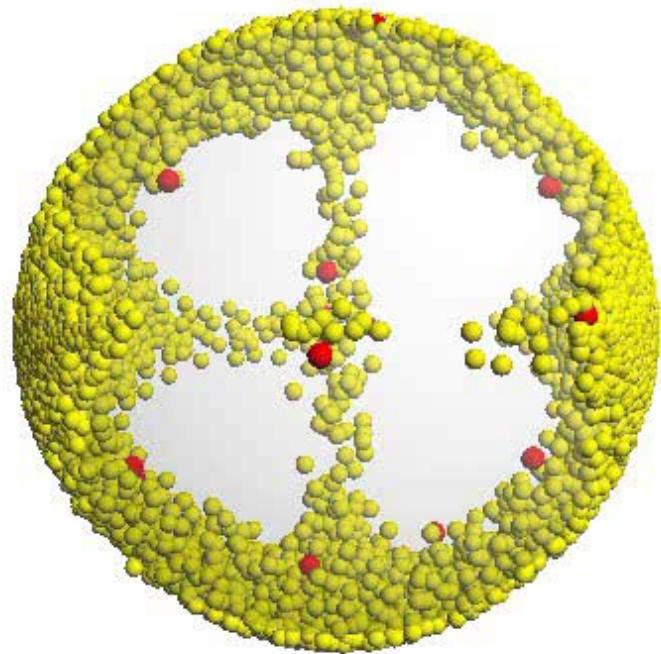


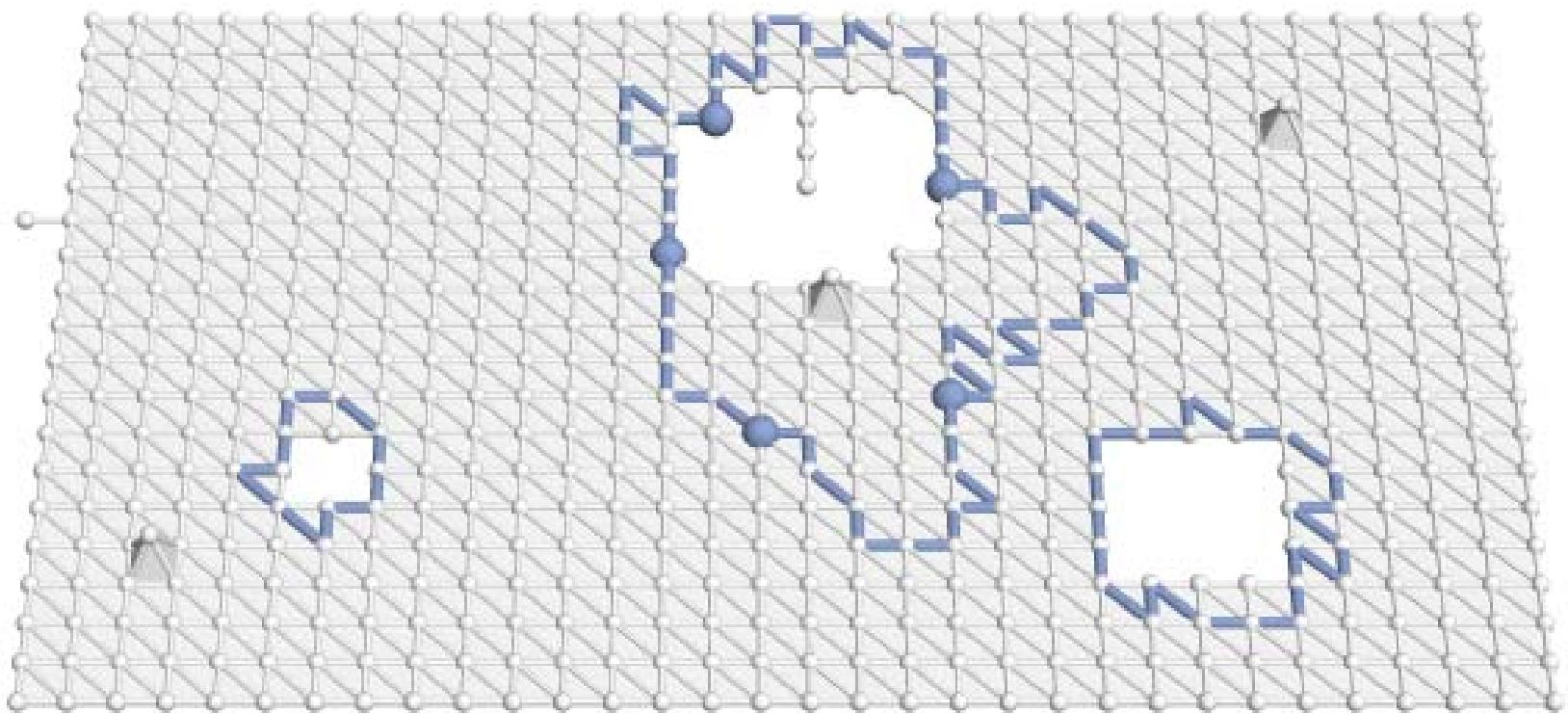


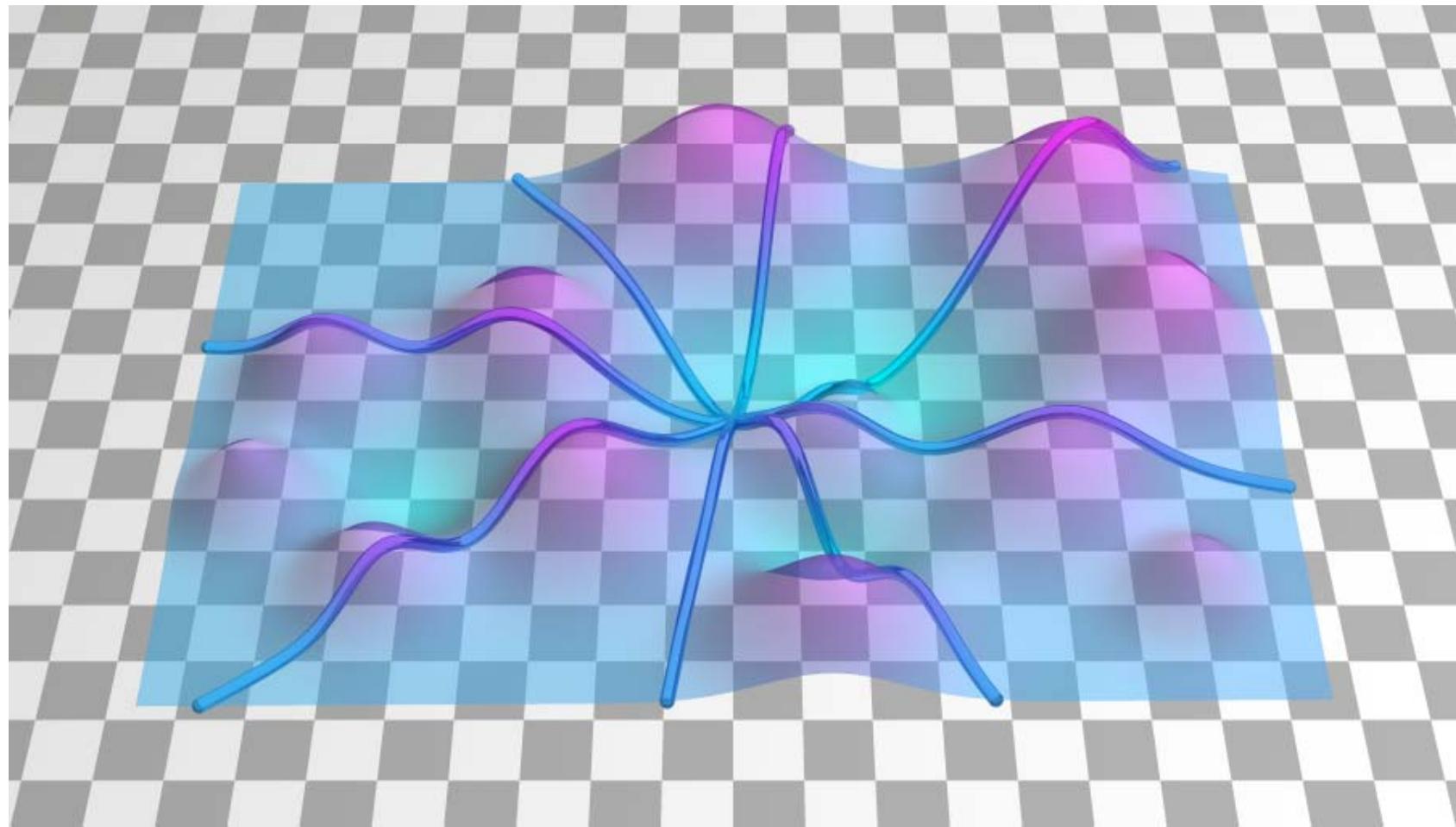






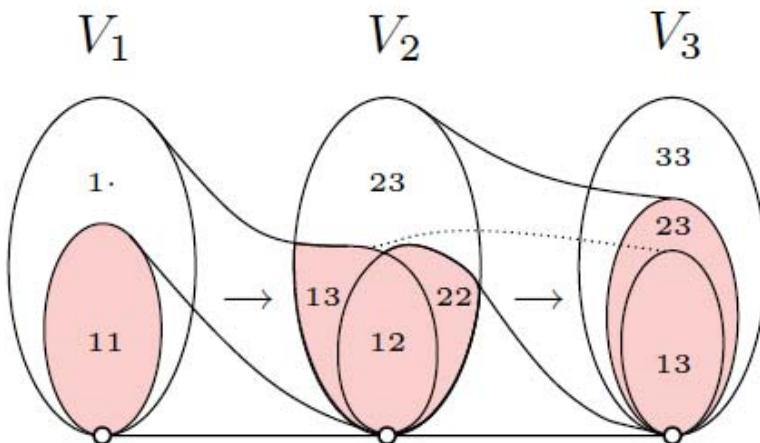




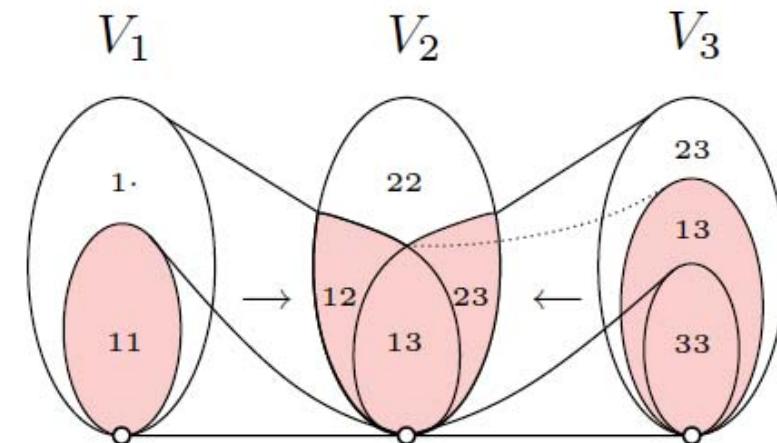


Zigzag Combinations

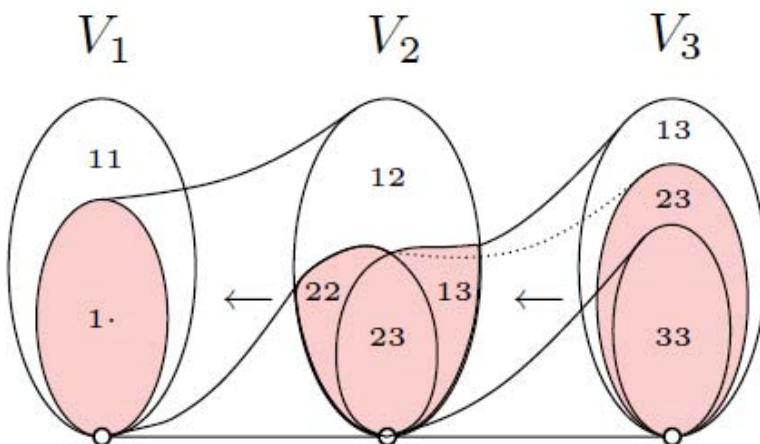
$V_1 \rightarrow V_2 \rightarrow V_3$



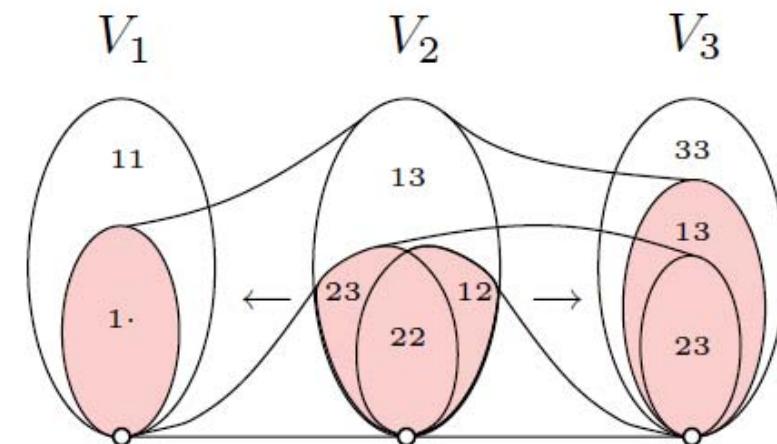
$V_1 \rightarrow V_2 \leftarrow V_3$



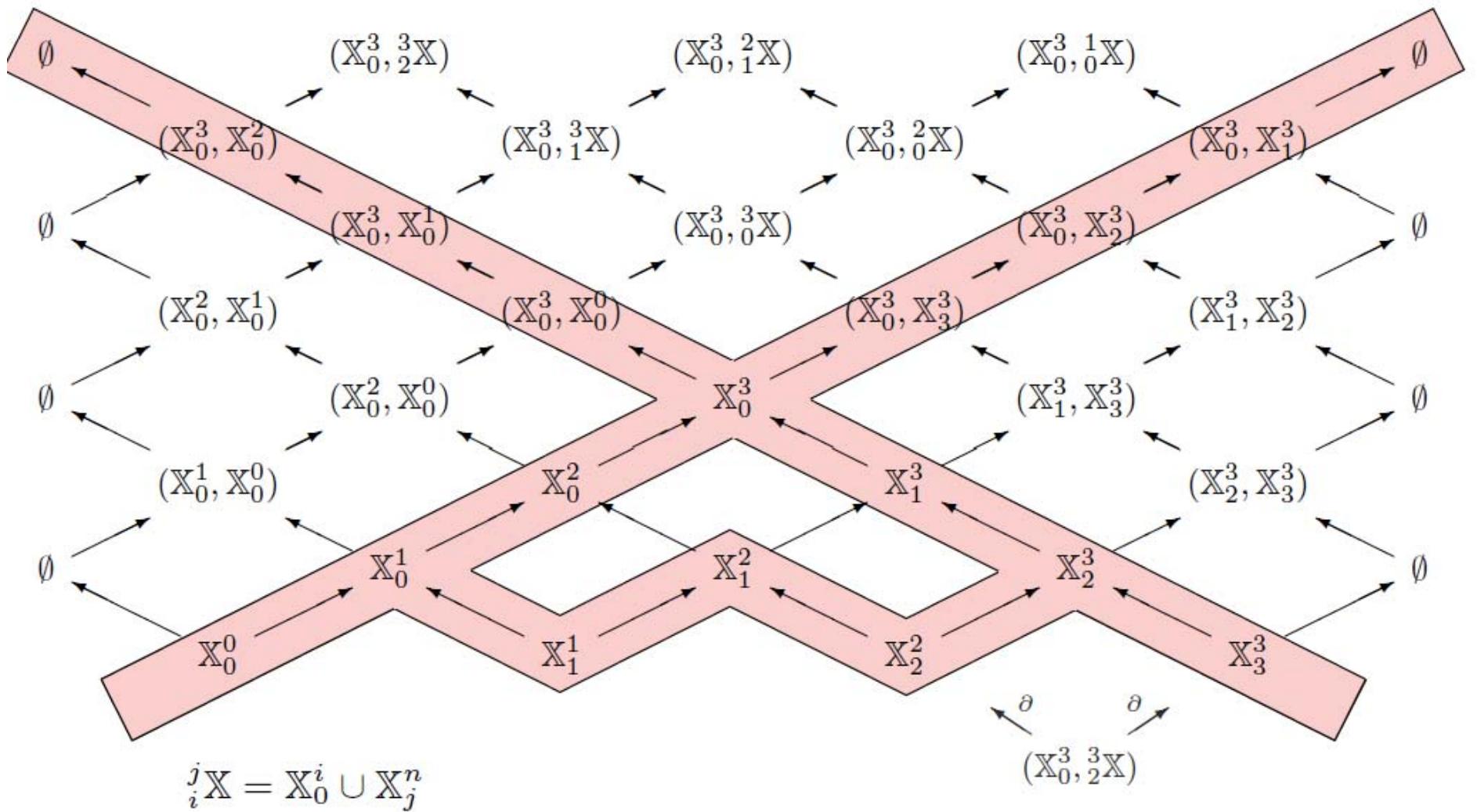
$V_1 \leftarrow V_2 \leftarrow V_3$



$V_1 \leftarrow V_2 \rightarrow V_3$



Pyramid



Approximating Shortest Homology Loops

Presenter: Tamal K. Dey

Department of CS&E, The Ohio State University

Joint work with Jian Sun and Yusu Wang



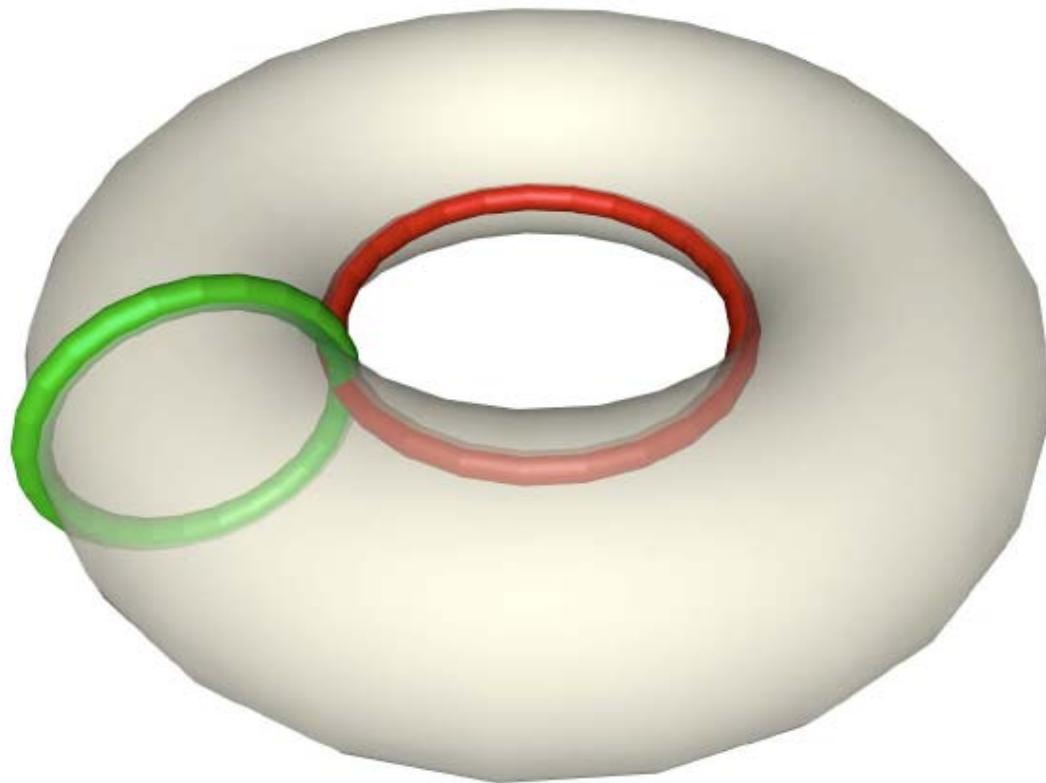
Abstract



- Inference of topology and geometry of a hidden manifold from its point data is a fundamental problem.
- We present an algorithm to compute a set of loops from point data that approximate a **shortest** basis of the homology group $H_1(M)$ of the sampled manifold M .
- We also present a polynomial time algorithm for computing a shortest basis of $H_1(M)$ for any finite simplicial complex embedded in an Euclidean space.



Example



Main Results: Theorem 1.3



- Let $M \subset \mathbb{R}^d$ be a smooth, closed manifold with l as the length of a shortest basis of $\mathbf{H}_1(M)$. Let $k = \text{rank } \mathbf{H}_1(M)$, $\rho(M)$ and $\rho_c(M)$ be *reach* and *convexity radius* of M , and $c(n, \alpha)$ be the worst-case size of any α -Rips complex with n vertices.
- Given a set $P \subset M$ of n points which is an ε -sample of M and $4\varepsilon \leq \alpha \leq \min\{\frac{1}{2}\sqrt{\frac{3}{5}}\rho(M), \rho_c(M)\}$, one can compute a set of loops G in $O(n^3 + knc(n, \alpha)^3)$ time where

$$\frac{1}{1 + \frac{4\alpha^2}{3\rho^2(M)}}l \leq \text{Len}(G) \leq \left(1 + \frac{4\varepsilon}{\alpha}\right)l.$$



Main Results: Theorem 1.4

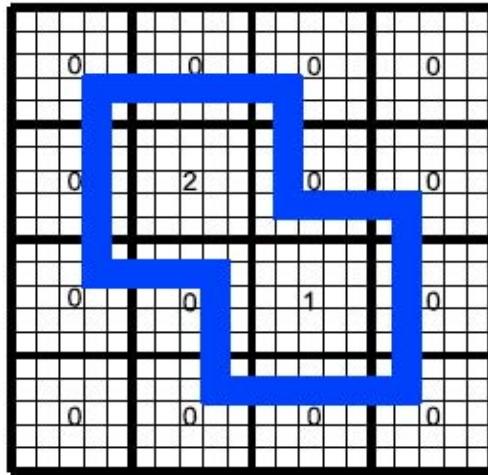


- Let \mathcal{K} be a finite simplicial complex with positive weights on edges.
- A shortest basis for $H_1(\mathcal{K})$ can be computed in $O(kn^4)$ time where $k = \text{rank } H_1(\mathcal{K})$ and $n = |\mathcal{K}|$.



Detection and Classification of Critical Components

Subdivide a cubical grid \mathcal{A} into five cubes in each dimension. Build an isolating neighbourhood N by taking two layers of subdivided cubes around $|\mathcal{A}|$. The outermost layer is the outer boundary, and will be divided into an upper level set L_p , a lower level set L_n and a level set L_z . The algorithm has two passes: first, cubes are identified as singular and consolidated into connected components; second, the components are identified as regular or critical, and classified.



Algorithm 1 (Detecting singular cubes)

```
For each elementary full cube Q
build N, Lp, Ln
H := H*(N, Lp) = 0 and H*(N, Ln) = 0
if H = TRUE then Q is ordinary
    else Q is singular
endif
```

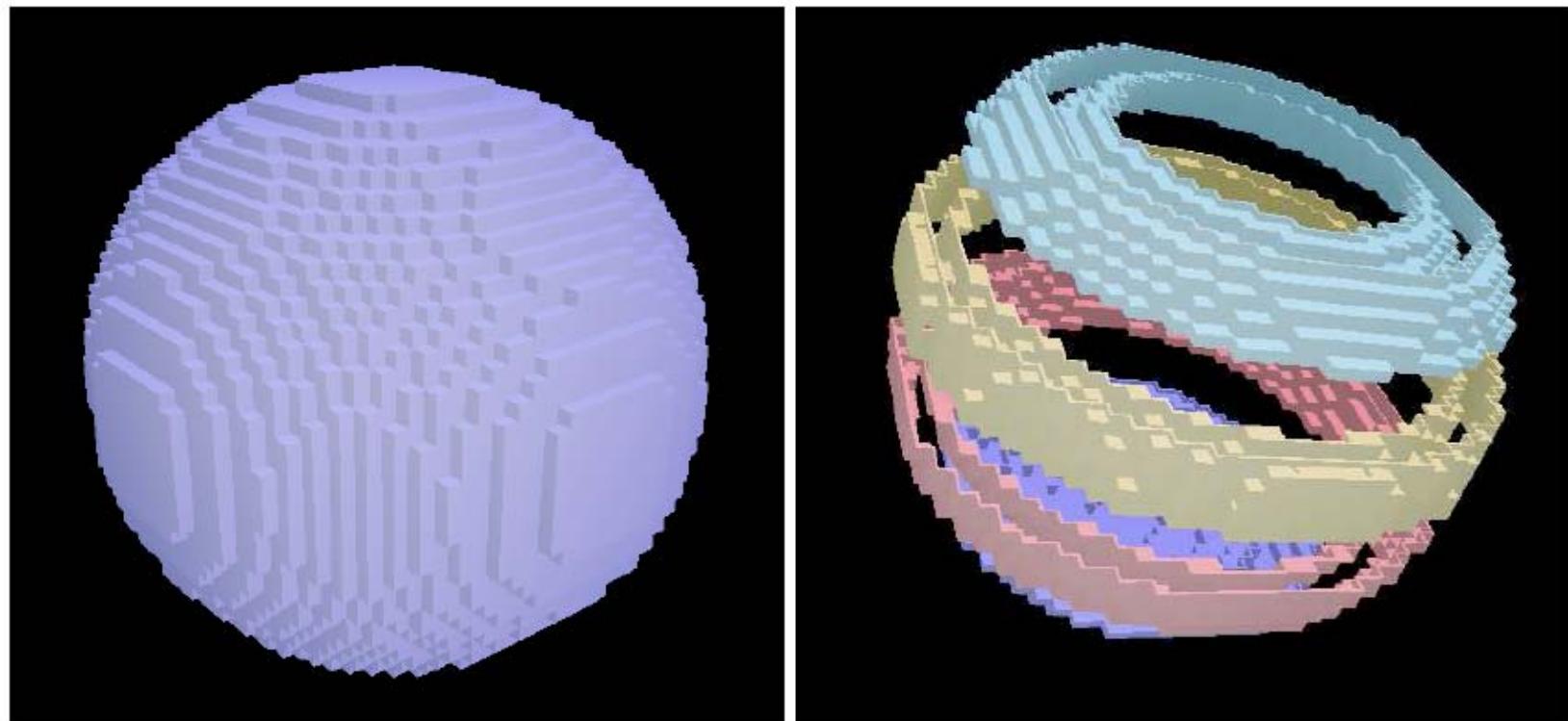
Algorithm 2 (Sorting components in \mathbb{R}^d)

```
For each singular cube Q, C := {Q}
while P ∈ bd(C) ∩ X is singular
    C := C ∪ {P}
    build bd(C)
endwhile
build N, Lp, Ln, Lz
do
    H := H*(N, Lp) = 0 and H*(N, Ln) = 0
    if H = TRUE then C is a regular component
    else if Lp = Lz = ∅
        then C is a maximum component
    else if Ln = Lz = ∅
        then C is a minimum component
    else C is a saddle component
    endif
break
```

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Critical Sphere



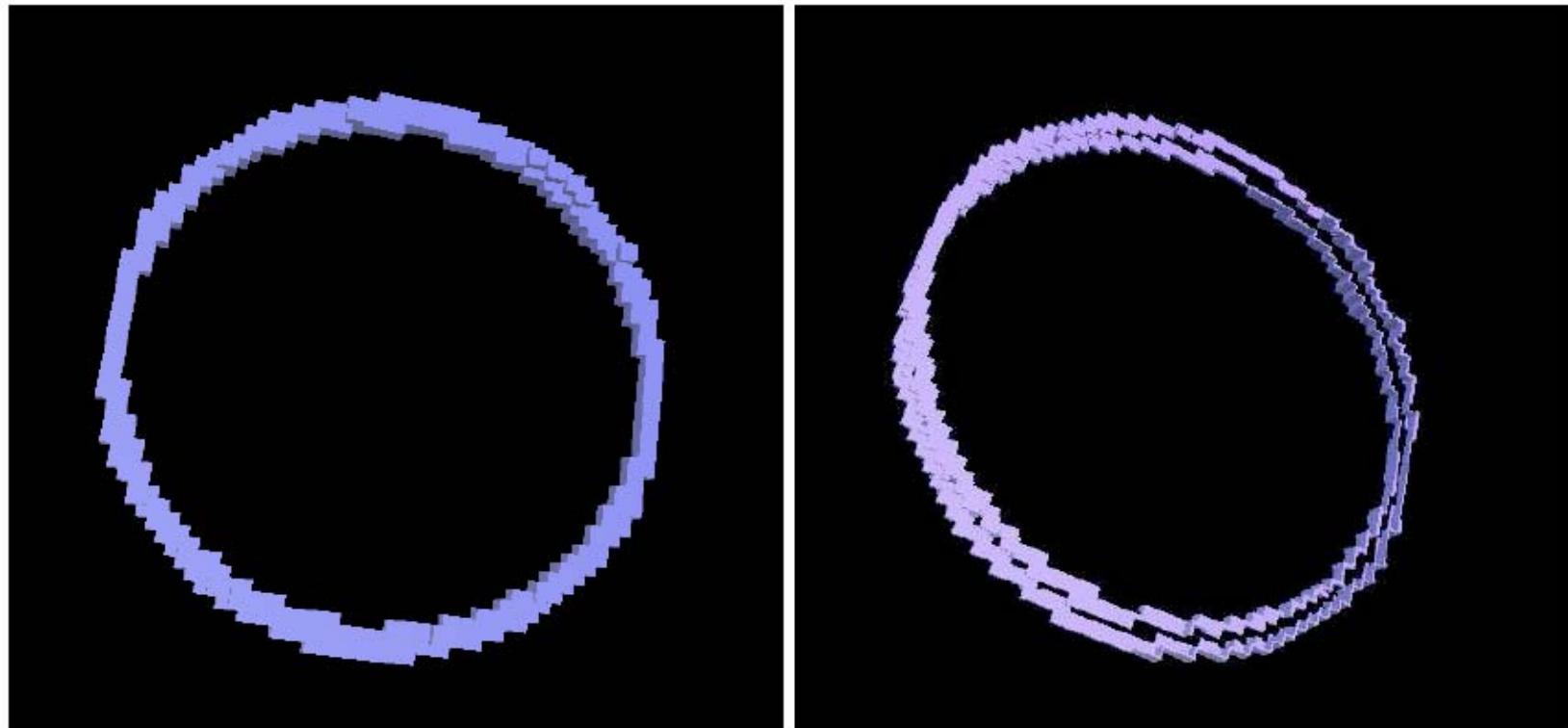
N and slices of L_p for the minimal sphere of $f(x, y, z) = (x^2 + y^2 + z^2 - 5)^2$

Domain: $[-3, 3]^3$ discretized in 51^3 3-cubes

Image: Discretized over 8192 levels

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Critical Torus



N and L_p for the function in toric coordinates $f(r, \theta, \phi) = r^2(1 - 2\sin^2(\theta - \phi/2))$

Domain: $[-3, 3]^2 \times [-4/5, 4/5]$ discretized in $101^2 \times 27$ 3-cubes

Image: Discretized over 16384 levels

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Madjid Allili, Bishop's University

