Building the Next Generation of Parallel Applications

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Quiz (True or False)

1. MPI-only has the best parallel performance.
2. Future parallel applications will not have MPI_Init().
3. All future programmers will need to write parallel code.
4. Use of “markup”, e.g., OpenMP pragmas, is the least intrusive approach to parallelizing a code.
5. DRY is not possible across CPUs and GPUs.
6. Extended precision is too expensive to be useful.
7. Resilience will be built into algorithms.
8. GPUs are a harbinger of CPU things to come.
9. Fortran Developers are in trouble in a manycore world.
10. Global SIMT is sufficient parallelism for scientific computing.
Trilinos Contributors

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Patrick Knupp
Joe Kotulska
Jason Kraftcheck
Trilinos Download History: 17600 Total
Registered User by Region

Registered Users by Region (5640 Total)

- Europe: 2105
- US (except Sandia): 1912
- Sandia (includes unregistered): 351
- Asia: 265
- Americas (except US): 866
- Australia/NZ: 65
- Africa: 76

Total Registered Users: 5640
Registered Users by Type
(5640 Total)

- University: 3369
- Government: 828
- Personal: 688
- Industry: 622
- Other: 133
Ubuntu/Debian: Other sources

maherou@jaguar13:/ccs/home/maherou> module avail trilinos

------------------------------------------------ /opt/
cray/modulefiles-------------------------------------------------
trilinos/10.0.1(default) trilinos/10.2.0

--------------------------------------------------- /sw/xt5/modulefiles--------------------------------------------------
trilinos/10.0.4 trilinos/10.2.2 trilinos/10.4.0 trilinos/8.0.3 trilinos/9.0.2
Capability Leaders: Layer of Proactive Leadership

- Areas:
  - Framework, Tools & Interfaces (J. Willenbring).
  - Software Engineering Technologies and Integration (R. Bartlett).
  - Discretizations (P. Bochev).
  - Scalable Linear Algebra (M. Heroux).
  - Linear & Eigen Solvers (J. Hu).
  - Nonlinear, Transient & Optimization Solvers (A. Salinger).
  - **Scalable I/O:** (R. Oldfield)

- Each leader provides strategic direction across all Trilinos packages within area.
# Trilinos Package Summary

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<th>Package(s)</th>
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<td>Teuchos, EpetraExt, <strong>Kokkos</strong>, Triutils, ThreadPool, Phalanx</td>
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<td>Solvers</td>
<td>AztecOO, Belos, Komplex</td>
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<td>Iterative linear solvers</td>
<td>Amesos, Amesos2</td>
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<tr>
<td>Direct sparse linear solvers</td>
<td>Epetra, Teuchos</td>
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<tr>
<td>Direct dense linear solvers</td>
<td>Anasazi, Rbgen</td>
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<tr>
<td>Iterative eigenvalue solvers</td>
<td>AztecOO, IFPACK, Ifpack2</td>
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<td>ILU-type preconditioners</td>
<td>ML, CLAPS</td>
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<td>Multilevel preconditioners</td>
<td>Meros, Teko</td>
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<td>Block preconditioners</td>
<td>NOX, LOCA</td>
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<td>Nonlinear system solvers</td>
<td>MOOCHO, Aristos, TriKota, Globipack, Optipack</td>
</tr>
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<td>Optimization (SAND)</td>
<td>Stochastic PDEs</td>
</tr>
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Three Design Points

• Terascale Laptop: Uninode-Manycore
• Petascale Deskside: Multinode-Manycore
• Exascale Center: Manynode-Manycore
Basic Concerns: Trends, Manycore

• Stein’s Law: *If a trend cannot continue, it will stop.*
  Herbert Stein, chairman of the Council of Economic Advisers under Nixon and Ford.

• Trends at risk:
  – Power.
  – Single core performance.
  – Node count.
  – Memory size & BW.
  – Concurrency expression in existing Programming Models.
  – Resilience.

![Parallel CG Performance 512 Threads](image)

32 Nodes = 2.2GHz AMD 4sockets X 4cores

“Status Quo” ~ MPI-only

Strong Scaling Potential

Edwards: SAND2009-8196
Trilinos ThreadPool Library v1.1.
Observations

• MPI-Only is not sufficient, except ... much of the time.
• Near-to-medium term:
  – MPI+[OMP|TBB|Pthreads|CUDA|OCL|MPI]
  – Long term, too?
• Concern:
  – Best hybrid performance: 1 MPI rank per UMA core set.
  – UMA core set size growing slowly ➔ Lots of MPI tasks.
• Long-term:
  – Something hierarchical, global in scope.
• Conjecture:
  – Data-intensive apps need non-SPDM model.
  – Will develop new programming model/env.
  – Rest of apps will adopt over time.
  – Time span: 10-20 years.
What Can we Do Right Now?

• Study why MPI was successful.
• Study new parallel landscape.
• Try to cultivate an approach similar to MPI (and others).
MPI Impresssions
MPI: It Hurts So Good

- Observations
  - “assembly language” of parallel computing
  - lowest common denominator
    - portable across architectures
  - upfront effort required
    - system design
    - application development

- Conclusions

So What Would Life Be Like Without MPI?

```
F(n) = \begin{cases} 
  1 & n = 0 \\
  0 & n = 1 \\
  F(n-1) + F(n-2) & n > 1 
\end{cases}
```

```
def fib(n):
    if n <= 1:
        return n
    else:
        return fib(n-1) + fib(n-2)
```

```
```

“MPI is often considered the “portable assembly language” of parallel computing, …”
3D Stencil in NAS MG

MPI Reality

How much MPI-specific code?
dft_fill_wjdc.c
MPI-specific code
MFIX: Multiphase Flows with Interphase eXchanges (https://www.mfix.org/)

**MFIX**

Source term for pressure correction

- MPI-callable, OpenMP-enabled.
- 340 Fortran lines.
- No MPI-specific code.
- Ubiquitous OpenMP markup (red regions).

source_pp_g.f
Reasons for MPI Success?

• Portability? Yes.
• Standardized? Yes.
• Momentum? Yes.
• Separation of many Parallel & Algorithms concerns? Big Yes.

• Once framework in place:
  – Sophisticated physics added as serial code.
  – Ratio of science experts vs. parallel experts: 10:1.

• Key goal for new parallel apps: Preserve this ratio
Computational Domain Expert Writing Future Parallel Code
Evolving Parallel Programming Model
Parallel Programming Model: Multi-level/Multi-device

Inter-node/inter-device (distributed) parallelism and resource management

Node-local control flow (serial)

Intra-node (manycore) parallelism and resource management

Stateless computational kernels run on each core

network of computational nodes

computational node with manycore CPUs and / or GPGPU

Message Passing

Threading

stateless kernels

Adapted from slide of H. Carter Edwards
Domain Scientist’s Parallel Palette

• MPI-only (SPMD) apps:
  – Single parallel construct.
  – Simultaneous execution.
  – Parallelism of even the messiest serial code.

• MapReduce:
  – Plug-n-Play data processing framework - 80% Google cycles.

• Pregel: Graph framework (other 20%)

• Next-generation PDE and related applications:
  – Internode:
    • MPI, yes, or something like it.
    • Composed with intranode.
  – Intranode:
    • Much richer palette.
    • More care required from programmer.

• What are the constructs in our new palette?
Obvious Constructs/Concerns

• Parallel for:
  – No loop-carried dependence.
  – Rich loops.
  – Use of shared memory for temporal reuse, efficient device data transfers.

• Parallel reduce:
  – Couple with other computations.
  – Concern for reproducibility.
Other construct: Pipeline

• Sequence of filters.
• Each filter is:
  – Sequential (grab element ID, enter global assembly) or
  – Parallel (fill element stiffness matrix).
• Filters executed in sequence.
• Programmer’s concern:
  – Determine (conceptually): Can filter execute in parallel?
  – Write filter (serial code).
  – Register it with the pipeline.
• Extensible:
  – New physics feature.
  – New filter added to pipeline.
TBB Pipeline for FE assembly

Launch elem-data from mesh
Serial Filter

Compute stiffnesses & loads
Parallel Filter

Assemble rows of stiffness into global matrix
Several Serial Filters in series

Each assembly filter assembles certain rows from a stiffness, then passes it on to the next assembly filter

Element-stiffness matrices computed in parallel

TBB work of Alan Williams
Alternative
TBB Pipeline for FE assembly

Launch elem-data from mesh

Compute stiffnesses & loads

Assemble rows of stiffness into global matrix

Serial Filter

Parallel Filter

Parallel Filter

Element-stiffness matrices computed in parallel

Each parallel call to the assembly filter assembles all rows from the stiffness, using locking to avoid race conflicts with other threads.
Base-line FE Assembly Timings

Problem size: 80x80x80 == 512000 elements, 531441 matrix-rows
The finite-element assembly performs 4096000 matrix-row sum-into operations (8 per element) and 4096000 vector-entry sum-into operations.

MPI-only, no threads. Linux dual quad-core workstation.

<table>
<thead>
<tr>
<th>Num-procs</th>
<th>Assembly-time Intel 11.1</th>
<th>Assembly-time GCC 4.4.4</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1.80s</td>
<td>1.95s</td>
</tr>
<tr>
<td>4</td>
<td>0.45s</td>
<td>0.50s</td>
</tr>
<tr>
<td>8</td>
<td>0.24s</td>
<td>0.28s</td>
</tr>
</tbody>
</table>
FE Assembly Timings

Problem size: 80x80x80 == 512000 elements, 531441 matrix-rows

The finite-element assembly performs 4096000 matrix-row sum-into operations (8 per element) and 4096000 vector-entry sum-into operations.

No MPI, only threads. Linux dual quad-core workstation.
Other construct: Thread team

• Multiple threads.
• Fast barrier.
• Shared, fast access memory pool.
• Example: Nvidia SM
• X86 more vague, emerging more clearly in future.
Observe: Iteration count increases with number of subdomains.

With scalable threaded smoothers (LU, ILU, Gauss-Seidel):
- Solve with fewer, larger subdomains.
- Better kernel scaling (threads vs. MPI processes).
- Better convergence, More robust

Exascale Potential: Tiled, pipelined implementation.

Three efforts:
- Level-scheduled triangular sweeps (ILU solve, Gauss-Seidel).
- Decomposition by partitioning
- Multithreaded direct factorization

Level Set Triangular Solver

Triangular Solve:
• Critical Kernel
  - MG Smoothers
  - Incomplete IC/ILU
• Naturally Sequential
• Building on classic algorithms:
  • Level Sched:
    • circa 1990.
  • Vectorization.
  • New: Generalized.

\[
\tilde{L} = PLPT = \begin{bmatrix}
D_1 & & \\
A_{2,1} & D_2 & \\
A_{3,1} & A_{3,2} & D_3 \\
& \vdots & \vdots & \ddots \\
A_{l,1} & A_{l,2} & A_{l,3} & \ldots & D_l
\end{bmatrix}
\]

\[
\tilde{x}_1 = D_1^{-1} \tilde{y}_1
\]
\[
\tilde{x}_2 = D_2^{-1} (\tilde{y}_2 - A_{2,1} \tilde{x}_1)
\]
\[
\vdots
\]
\[
\tilde{x}_l = D_l^{-1} (\tilde{y}_l - A_{l,1} \tilde{x}_1 - \ldots - A_{l,l-1} \tilde{x}_{l-1})
\]

Permutated System

Multi-step Algorithm
Triangular Solve Results

Passive (PB) vs. Active (AB) Barriers: Critical for Performance

<table>
<thead>
<tr>
<th>Name</th>
<th>N</th>
<th>nnz</th>
<th>N/nlevels</th>
<th>application area</th>
</tr>
</thead>
<tbody>
<tr>
<td>asic680ks</td>
<td>682,712</td>
<td>2,329,176</td>
<td>13932.9</td>
<td>circuit simulation</td>
</tr>
<tr>
<td>cage12</td>
<td>130,228</td>
<td>2,032,536</td>
<td>1973.2</td>
<td>DNA electrophoresis</td>
</tr>
<tr>
<td>pkustk04</td>
<td>55,590</td>
<td>4,218,660</td>
<td>149.4</td>
<td>structural engineering</td>
</tr>
<tr>
<td>bcsstk32</td>
<td>44,609</td>
<td>2,014,701</td>
<td>15.1</td>
<td>structural engineering</td>
</tr>
</tbody>
</table>

Level sets: Trilinos/Isorropia
Core Kernel Timings: Trilinos/Kokkos.
Thread Team Advantages

- Qualitatively better algorithm:
  - Threaded triangular solve scales.
  - Fewer MPI ranks means fewer iterations, better robustness.

- Exploits:
  - Shared data.
  - Fast barrier.
  - Data-driven parallelism.
Finite Elements/Volumes/Differences and parallel node constructs

- Parallel for, reduce, pipeline:
  - Sufficient for vast majority of node level computation.
  - Supports:
    - Complex modeling expression.
    - Vanilla parallelism.
  - Must be “Stencil-aware” for temporal locality.

- Thread team:
  - Complicated.
  - Requires true parallel algorithm knowledge.
  - Useful in solvers.
Portable Multi/Manycore Programming
Trilinos/Kokkos Node API
Another Manycore architecture: Intel MIC

Knights Ferry:

- 32 x86 cores
  - 4-way hyperthreading
  - 128 threads total
- 512-bit vector unit
  - 16 floats, 8 doubles
- 1.20GHz
- PCI-E 2.0
- 2GB GDDR5 global mem
  - 8MB shared L2 cache
  - 64KB L1 Data, 64KB L1 Inst

Programming Env:

- OpenMP,
- TBB,
- Pthreads
Tpetra and Kokkos

• **Tpetra** is an implementation of the Petra Object Model.
  – Design is similar to Epetra, with appropriate deviation.
  – Fundamental differences:
    • heavily exploits templates
    • utilizes hybrid (distributed + **shared**) parallelism via Kokkos Node API

• **Kokkos** is an API for shared-memory parallel nodes
  – Provides parallel_for and parallel_reduce skeletons.
  – Support shared memory APIs:
    • ThreadPool Interface (TPI; Carter Edwards’s pthreads Trilinos package)
    • Intel Threading Building Blocks (TBB)
    • NVIDIA CUDA-capable GPUs (via Thrust)
    • **OpenMP** (implemented by Radu Popescu/EPFL, awaiting my git push)
Generic Shared Memory Node

• Abstract inter-node comm provides DMP support.
• Need some way to *portably* handle SMP support.
• Goal: allow code, once written, to be run on *any parallel node*, regardless of architecture.

• **Difficulty #1**: Many different memory architectures
  – Node may have multiple, disjoint memory spaces.
  – Optimal performance may require special memory placement.

• **Difficulty #2**: Kernels must be tailored to architecture
  – Implementation of optimal kernel will vary between archs
  – No universal binary $\Rightarrow$ need for separate compilation paths

• Practical goal: Cover 80% kernels with generic code.
Kokkos Node API

• **Kokkos** provides two main components:
  – **Kokkos memory model** addresses Difficulty #1
    • Allocation, deallocation and efficient access of memory
    • **compute buffer**: special memory used for parallel computation
    • New: Local Store Pointer and Buffer with size.
  – **Kokkos compute model** addresses Difficulty #2
    • Description of kernels for parallel execution on a node
    • Provides stubs for common parallel work constructs
    • Currently, **parallel for loop** and **parallel reduce**

• Code is developed around a polymorphic Node object.
• Supporting a new platform requires only the implementation of a new node type.
Kokkos Memory Model

• A generic node model must at least:
  – support the scenario involving distinct device memory
  – allow efficient memory access under traditional scenarios
• Nodes provide the following memory routines:

  ArrayRCP<T> Node::allocBuffer<T>(size_t sz);
  void       Node::copyToBuffer<T>(  T * src,
                                      ArrayRCP<T>  dest);
  void       Node::copyFromBuffer<T>(ArrayRCP<T> src,
                                      T * dest);
  ArrayRCP<T> Node::viewBuffer<T> (ArrayRCP<T> buff);
  void       Node::readyBuffer<T>(ArrayRCP<T> buff);
Kokkos Compute Model

- How to make shared-memory programming generic:
  - Parallel reduction is the intersection of \texttt{dot()} and \texttt{norm1()}
  - Parallel for loop is the intersection of \texttt{axpy()} and mat-vec
  - We need a way of \texttt{fusing} kernels with these basic \texttt{constructs}.
- Template meta-programming is \texttt{the answer}.
  - This is the same approach that Intel TBB and Thrust take.
  - Has the effect of requiring that Tpetra objects be templated on Node type.
- Node provides generic parallel constructs, user fills in the rest:

```cpp
// Work-data pair (WDP) struct provides:
//  - loop body via \texttt{WDP::execute(i)}

template <class WDP>
void Node::\texttt{parallel_for}(
  int beg, int end, WDP workdata);

// Work-data pair (WDP) struct provides:
//  - reduction type \texttt{WDP::ReductionType}
//  - element generation via \texttt{WDP::generate(i)}
//  - reduction via \texttt{WDP::reduce(x,y)}

template <class WDP>
WDP::ReductionType Node::\texttt{parallel_reduce}(
  int beg, int end, WDP workdata);
```
Example Kernels: \texttt{axpy()} and \texttt{dot()}

\begin{verbatim}
template <class WDP>
void
Node::parallel_for(int beg, int end, WDP workdata);

template <class WDP>
WDP::ReductionType
Node::parallel_reduce(int beg, int end, WDP workdata);

template <class T>
struct AxpyOp {
    const T * x;
    T * y;
    T alpha, beta;
    void execute(int i)
    { y[i] = alpha*x[i] + beta*y[i]; }
};

AxpyOp<double> op;
op.x = ...;  op.alpha = ...;
op.y = ...;  op.beta = ...;
node.parallel_for< AxpyOp<double> > (0, length, op);

template <class T>
struct DotOp {
    typedef T ReductionType;
    const T * x, * y;
    T identity() { return (T)0; }
    T generate(int i) { return x[i]*y[i]; }
    T reduce(T x, T y) { return x + y; }
};

DotOp<float> op;
op.x = ...;  op.y = ...;
float dot;
dot = node.parallel_reduce< DotOp<float> > (0, length, op);
\end{verbatim}
What’s the Big Deal about Vector-Vector Operations?

Examples from OOQP (Gertz, Wright)

\[ y_i \leftarrow y_i + \alpha x_i z_i \quad ,i = 1...n \]
\[ y_i \leftarrow \begin{cases} y_{\text{min}} - y_i & \text{if } y_i < y_{\text{min}} \\ y_{\text{max}} - y_i & \text{if } y_i > y_{\text{max}} \\ 0 & \text{if } y_{\text{min}} \leq y_i \leq y_{\text{max}} \end{cases} ,i = 1...n \]
\[ y_i \leftarrow y_i / x_i \quad ,i = 1...n \]
\[ \alpha \leftarrow \{ \text{max } \alpha : x + \alpha d \geq \beta \} \]

Example from TRICE (Dennis, Heinkenschloss, Vicente)

\[ d_i \leftarrow \begin{cases} (b - u)_i^{1/2} & \text{if } w_i < 0 \text{ and } b_i < +\infty \\ 1 & \text{if } w_i < 0 \text{ and } b_i = +\infty \\ (u - a)_i^{1/2} & \text{if } w_i \geq 0 \text{ and } a_i > -\infty \\ 1 & \text{if } w_i \geq 0 \text{ and } a_i = -\infty \end{cases} ,i = 1...n \]

Example from IPOPT (Waechter)

\[ x_i \leftarrow \begin{cases} x_i^L + \frac{x_i^U - x_i^L}{2} & \text{if } \hat{x}_i^L > \hat{x}_i^U \\ \hat{x}_i^L & \text{if } x_i < \hat{x}_i^L \quad ,i = 1...n \\ \hat{x}_i^U & \text{if } x_i > \hat{x}_i^U \end{cases} \]

where:

\[ \hat{x}_i^L = \min \left( x_i^L + \eta \left( x_i^U - x_i^L \right) x_i^L + \delta \right) \]
\[ \hat{x}_i^U = \max \left( x_i^L - \eta \left( x_i^U - x_i^L \right) x_i^U - \delta \right) \]

Many different and unusual vector operations are needed by interior point methods for optimization!

Currently in MOOCHO:

- > 40 vector operations!
Tpetra RTI Components

• Set of stand-alone non-member methods:
  – unary_transform<UOP>(Vector &v, UOP op)
  – binary_transform<BOP>(Vector &v1, const Vector &v2, BOP op)
  – reduce<G>(const Vector &v1, const Vector &v2, G op_glob)
  – binary_pre_transform_reduce<G>(Vector &v1,
      const Vector &v2,
      G op_glob)

• These are non-member methods of Tpetra::RTI which are loosely coupled with Tpetra::MultiVector and Tpetra::Vector.

• Tpetra::RTI also provides Operator-wrappers:
  – class KernelOp<..., Kernel > : Tpetra::Operator<...>
  – class BinaryOp<...,BinaryOp> : Tpetra::Operator<...>
Tpetra RTI Example

// isn’t this nicer than a bunch of typedefs?
auto &platform = Tpetra::DefaultPlatform::getDefaultPlatform();
auto comm = platform.getComm();
auto node = platform.getNode();

// create Map and some Vector objects
Tpetra::global_size_t numGlobalRows = ...;
auto map = createUniformContigMapWithNode<int,int>(numGlobalRows, comm, node);
const size_t numLocalRows = map->getNodeNumElements();
auto x = Tpetra::createVector<float>(map),
      y = Tpetra::createVector<float>(map);
auto z = Tpetra::createVector<double>(map),
      w = Tpetra::createVector<double>(map);

// parallel initialization of x[i] = 1.0 using C++-0x lambda function
Tpetra::RTI::unary_transform( *x, [](float xi){return 1.0f;} );
// parallel initialization of y[i] = x[i]
Tpetra::RTI::binary_transform( *y, *x, [](float, float xi) {return xi;} );
// parallel y[i] = x[i] + y[i]
Tpetra::RTI::binary_transform( *y, *x, std::plus<float>() );
// parallel single precision dot(x,y)
float result = Tpetra::RTI::reduce( *x, *y, reductionGlob<ZeroOp<float>>(
                                             std::multiplies<float>(),
                                             std::plus<float>() ));
**Multiprecision possibilities**

- **Tpetra** is a templated version of the Petra distributed linear algebra model in Trilinos.
  - Objects are templated on the underlying data types:
    ```
    MultiVector<scalar=double, local_ordinal=int, global_ordinal=local_ordinal> ...
    CrsMatrix<scalar=double, local_ordinal=int, global_ordinal=local_ordinal> ...
    ```
  - Examples:
    ```
    MultiVector<double, int, long int> V;
    CrsMatrix<float> A;
    ```

- Speedup of float over double in Belos linear solver.
  - |                | float | double | speedup |
  - |----------------|-------|--------|--------|
  - | Solve time (s) | 2.6   | 5.3    | 1.42x  |
  - | Accuracy       | \(10^{-6}\) | \(10^{-12}\) | \(10^{-24}\) | \(10^{-48}\) |

Arbitrary precision solves using Tpetra and Belos linear solver package.
FP Accuracy Analysis:
FloatShadowDouble Datatype

class FloatShadowDouble {
public:
    FloatShadowDouble() {
        f = 0.0f;
        d = 0.0;
    }
    FloatShadowDouble(const FloatShadowDouble &fd) {
        f = fd.f;
        d = fd.d;
    }
    ...
    inline FloatShadowDouble operator+=(const FloatShadowDouble &fd) {
        f += fd.f;
        d += fd.d;
        return *this;
    }
    ...
    inline std::ostream& operator<<(std::ostream& os, const FloatShadowDouble& fd) {
        os << fd.f << "f" << fd.d << "d"; return os;
    }

- Templates enable new analysis capabilities
- Example: Float with “shadow” double.
## FloatShadowDouble

Sample usage:
```
#include "FloatShadowDouble.hpp"
Tpetra::Vector<FloatShadowDouble> x, y;
Tpetra::CrsMatrix<FloatShadowDouble> A;
A.apply(x, y);  // Single precision, but double results also computed, available
```

<table>
<thead>
<tr>
<th>Iteration</th>
<th>Residual</th>
<th>Residual</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>455.194f</td>
<td>455.194d</td>
</tr>
<tr>
<td>15</td>
<td>5.07328f</td>
<td>5.07618d</td>
</tr>
<tr>
<td>30</td>
<td>0.00147022f</td>
<td>0.00138466d</td>
</tr>
<tr>
<td>45</td>
<td>5.14891e-06f</td>
<td>2.09624e-06d</td>
</tr>
<tr>
<td>60</td>
<td>4.03386e-09f</td>
<td>7.91927e-10d</td>
</tr>
</tbody>
</table>
Hybrid CPU/GPU Computing
Writing and Launching Heterogeneous Jobs

• A node is a shared-memory domain.
• Multiple nodes are coupled via a communicator.
  – This requires launching multiple processes.
• In a heterogeneous cluster, this requires code written for multiple node types.
• It may be necessary to template large parts of the code and run the appropriate instantiation on each rank.
• For launching, two options are available:
  – Multiple single-node executables, complex dispatch
  – One diverse executable, early branch according to rank
Tpetra::HybridPlatform

• Encapsulate main in a templated class method:

```cpp
template <class Node>
class myMainRoutine {
    static void run(ParameterList &runParams,
                     const RCP<const Comm<int> > &comm,
                     const RCP<Node> &node)
    {
        // do something interesting
    }
};
```

• **HybridPlatform** maps the communicator rank to the Node type, instantiates a node and the run routine:

```cpp
int main(...) {
    Comm<int> comm = ...,
    ParameterList machine_file = ...
    // instantiate appropriate node and myMainRoutine
    Tpetra::HybridPlatform platform( comm , machine_file );
    platform.runUserCode< myMainRoutine >();
    return 0;
}
```
## HybridPlatform Machine File

<table>
<thead>
<tr>
<th>round-robin assignment</th>
<th>interval assignment</th>
<th>explicit assignment</th>
<th>default</th>
</tr>
</thead>
<tbody>
<tr>
<td>%M=N</td>
<td>[M,N]</td>
<td>=N</td>
<td>default</td>
</tr>
</tbody>
</table>

### Parameter List

```
<ParameterList>
  <ParameterList name="%2=0">
    <Parameter name="NodeType" type="string" value="Kokkos::ThrustGPUNode"/>
    <Parameter name="Verbose" type="int" value="1"/>
    <Parameter name="Device Number" type="int" value="0"/>
    <Parameter name="Node Weight" type="int" value="4"/>
  </ParameterList>
  <ParameterList name="%2=1">
    <Parameter name="NodeType" type="string" value="Kokkos::TPINode"/>
    <Parameter name="Verbose" type="int" value="1"/>
    <Parameter name="Num Threads" type="int" value="15"/>
    <Parameter name="Node Weight" type="int" value="15"/>
  </ParameterList>
</ParameterList>
```
HybridPlatformTest Output

[tpetra/example/HybridPlatform] mpirun -np 4 ./Tpetra_HybridPlatformTest.exe
--machine-file=machines/G+15.xml

Every proc machine parameters from: machines/G+15.xml

Teuchos::GlobalMPISession::GlobalMPISession(): started with name lens31 and rank 0!
Running test with Node == Kokkos::ThrustGPUNode on rank 0/4
ThrustGPUNode attached to device #0 "Tesla C1060", of compute capability 1.3

Teuchos::GlobalMPISession::GlobalMPISession(): started with name lens31 and rank 1!
Running test with Node == Kokkos::TPINode on rank 1/4

Teuchos::GlobalMPISession::GlobalMPISession(): started with name lens10 and rank 2!
Running test with Node == Kokkos::ThrustGPUNode on rank 2/4
TPINode initializing with numThreads == 15
ThrustGPUNode attached to device #0 "Tesla C1060", of compute capability 1.3

Teuchos::GlobalMPISession::GlobalMPISession(): started with name lens10 and rank 3!
Running test with Node == Kokkos::TPINode on rank 3/4
TPINode initializing with numThreads == 15

...

See HybridPlatformAnasazi.cpp and HybridPlatformBelos.cpp for more fun!
Programming Today for Tomorrow’s Machines
Programming Today for Tomorrow’s Machines

• Parallel Programming in the small:
  – Focus: writing sequential code fragments.
  – Programmer skills:
    • 10%: Pattern/framework experts (domain-aware).
    • 90%: Domain experts (pattern-aware)

• Languages needed are already here.
  – Exception: Large-scale data-intensive graph?
FE/FV/FD Parallel Programming Today

for ((i,j,k) in points/elements on subdomain) {
    compute coefficients for point (i,j,k)
    inject into global matrix
}

Notes:

- User in charge of:
  - Writing physics code.
  - Iteration space traversal.
  - Storage association.

- Pattern/framework/runtime in charge of:
  - SPMD execution.
FE/FV/FD Parallel Programming Tomorrow

pipeline \langle i,j,k \rangle \{
    filter(\text{addPhysicsLayer1}<i,j,k>);
    ...
    filter(\text{addPhysicsLayern}<i,j,k>);
    filter(\text{injectIntoGlobalMatrix}<i,j,k>);
}\}

Notes:

• User in charge of:
  – Writing physics code (filter).
  – Registering filter with framework.

• Pattern/framework/runtime in charge of:
  – SPMD execution.
  – Iteration space traversal.
    – Sensitive to temporal locality.
  – Filter execution scheduling.
  – Storage association.

• Better assignment of responsibility (in general).
Software Development and Delivery
Compile-time Polymorphism
Templates and Sanity upon a shifting foundation

Software delivery:
• Essential Activity

How can we:
• Implement mixed precision algorithms?
• Implement generic fine-grain parallelism?
• Support hybrid CPU/GPU computations?
• Support extended precision?
• Explore redundant computations?
• Prepare for both exascale “swim lanes”?

C++ templates only sane way:
• Moving to completely templated Trilinos libraries.
• Other important benefits.
• A usable stack exists now in Trilinos.

“Are C++ templates safe? No, but they are good.”

Template Benefits:
– Compile time polymorphism.
– True generic programming.
– No runtime performance hit.
– Strong typing for mixed precision.
– Support for extended precision.
– Many more…

Template Drawbacks:
– Huge compile-time performance hit:
  • But good use of multicore :)  
  • Eliminated for common data types.
– Complex notation:
  – Esp. for Fortran & C programmers).
  – Can insulate to some extent.
# Solver Software Stack

**Phase I packages: SPMD, int/double**  
**Phase II packages: Templated**

| **Optimization** | **Find** \( u \in \mathbb{R}^n \) that minimizes \( g(u) \)  
Unconstrained:  
Constrained: \( \text{Find} \ x \in \mathbb{R}^m \) and \( u \in \mathbb{R}^n \) that minimizes \( g(x, u) \) s.t. \( f(x, u) = 0 \) |
| **Bifurcation Analysis** | **Given nonlinear operator** \( F(x, u) \in \mathbb{R}^{n+m} \) \( \text{For} \ F(x, u) = 0 \) find space \( u \in U \oplus \frac{\partial F}{\partial x} \) |
| **Transient Problems** | **Solve** \( f(\dot{x}(t), x(t), t) = 0 \)  
DAEs/ODEs: \( t \in [0, T], x(0) = x_0, \dot{x}(0) = x'_0 \)  
for \( x(t) \in \mathbb{R}^n, t \in [0, T] \) |
| **Nonlinear Problems** | **Given nonlinear operator** \( F(x) \in \mathbb{R}^m \rightarrow \mathbb{R}^m \) \( \text{Solve} \ F(x) = 0 \) \( x \in \mathbb{R}^n \) |
| **Linear Problems** | **Given Linear Ops (Matrices)** \( A, B \in \mathbb{R}^{m \times n} \)  
Linear Equations: \( \text{Solve} \ Ax = b \) \( \text{for} \ x \in \mathbb{R}^n \)  
Eigen Problems: \( \text{Solve} \ A\nu = \lambda B\nu \) \( \text{for} \) (all) \( \nu \in \mathbb{R}^n, \lambda \in \mathbb{R} \) |
| **Distributed Linear Algebra** | **Compute** \( y = Ax; A = A(G); A \in \mathbb{R}^{m \times n}, G \in \mathbb{S}^{m \times n} \)  
Matrix/Graph Equations: \( \text{Compute} \ y = \alpha x + \beta w; \alpha = \langle x, y \rangle; x, y \in \mathbb{R}^n \)  
Vector Problems: |
### Solver Software Stack

<table>
<thead>
<tr>
<th>Phase I packages</th>
<th>Phase II packages</th>
<th>Phase III packages: Manycore*, templated</th>
</tr>
</thead>
</table>

#### Optimization
- **Unconstrained:** Find $u \in \mathbb{R}^n$ that minimizes $g(u)$
- **Constrained:** Find $x \in \mathbb{R}^m$ and $u \in \mathbb{R}^n$ that minimizes $g(x, u)$ s.t. $f(x, u) = 0$

#### Bifurcation Analysis
- Given nonlinear operator $F(x, u) \in \mathbb{R}^{n+m}$
- For $F(x, u) = 0$ find space $u \in U \ni \frac{\partial F}{\partial x}$

#### Transient Problems
- **DAEs/ODEs:** Solve $f(\dot{x}(t), x(t), t) = 0$
  - $t \in [0, T], x(0) = x_0, \dot{x}(0) = x_0'$
  - for $x(t) \in \mathbb{R}^n, t \in [0, T]$

#### Nonlinear Problems
- Given nonlinear operator $F(x) \in \mathbb{R}^m \rightarrow \mathbb{R}$
- Solve $F(x) = 0 \; x \in \mathbb{R}^n$

#### Linear Problems
- **Linear Equations:** Given Linear Ops (Matrices) $A, B \in \mathbb{R}^{m \times n}$
  - Solve $Ax = b \; x \in \mathbb{R}^n$
- **Eigen Problems:** Solve $Av = \lambda Bv$ for (all) $v \in \mathbb{R}^n, \lambda \in \mathbb{R}$

#### Distributed Linear Algebra
- **Matrix/Graph Equations:** Compute $y = Ax; A = A(G); A \in \mathbb{R}^{m \times n}, G \in \mathbb{S}^{m \times n}$
- **Vector Problems:** Compute $y = \alpha x + \beta w; \alpha = (x, y); x, y \in \mathbb{R}^n$

---

### Sensitivities

(Automatic Differentiation: Sacado)

- **MOOCHO**
- **LOCA**
- **T-LOCA**
- **Rythmos**
- **NOX**
- **T-NOX**
- **Anasazi**
- **AztecOO**
- **Belos***
- **Ifpack, ML, etc...**
- **T-Ifpack***
- **T-ML***, etc...
- **Epeta**
- **Tpetra***
- **Kokkos***
- **Teuchos**
C++: Great story

- Great native language support: TBB, CUDA.
- Kokkos: Generic interface layer.
- Trilinos built on top.
The Challenges for Fortran
Fortran Manycore Programming Environments

- **OpenMP**:
  - Only works on CPUs.
  - Some research efforts to migrate it to GPUs.
- **CUDA Fortran**:
  - Isomorphic to CUDA.
  - Proprietary.
  - No apparent effort to standardize syntax.
- **Multi-language**:
  - Use C/C++ for kernels.
  - Problem:
    - Every kernel must be manycore executable.
    - Implies a complete rewrite.
- Other option: A complete rewrite.
Underlying Challenges

• Desire:
  – Write the loop body only.
  – Hoist it up (at compile time, with runtime support) into the threaded execution environment.
  – CUDA: <<< >>> syntax.
  – TBB: parallel_for syntax.
  – OpenMP: Deconstruct loop for parallel execution.

• Problem:
  – No compile-time polymorphism.
  – OpenMP only works with
    • parallel_for, parallel_reduce, basic tasks.
    • No support for pipelines (yet).
Bi-Modal: MPI-only and MPI+\([X|Y|Z]\)
Parallel machine with \( p = m \times n \) processors:
- \( m = \) number of nodes.
- \( n = \) number of shared memory processors per node.

Two ways to program:
- Way 1: \( p \) MPI processes.
- Way 2: \( m \) MPI processes with \( n \) threads per MPI process.

New third way:
- “Way 1” in some parts of the execution (the app).
- “Way 2” in others (the solver).
Multicore Scaling: App vs. Solver

Application:
- Scales well
  (sometimes superlinear)
- MPI-only sufficient.

Solver:
- Scales more poorly.
- Memory system-limited.
- MPI+threads can help.

* Charon Results:
  Lin & Shadid TLCC Report
MPI-Only + MPI/Threading: $Ax=b$

- Application (App) passes matrix and vector values to library data classes.
- All ranks store $A$, $x$, $b$ data in memory visible to rank 0.
- Library solves $Ax=b$ using shared memory algorithms on the node.

**Multicore: “PNAS” Layout**

- Application (App) ranks 0 to 3.
- Library (Lib) ranks 0 to 3.
- Memory (Mem) ranks 0 to 3.
- Threads: Thread 0 to 3.
MPI Shared Memory Allocation

Idea:
- Shared memory alloc/free functions:
  - MPI_Comm_alloc_mem
  - MPI_Comm_free_mem
- Predefined communicators:
  - MPI_COMM_NODE – ranks on node
  - MPI_COMM_SOCKET – UMA ranks
  - MPI_COMM_NETWORK – inter node
- Status:
  - Available in current development branch of OpenMPI.
  - First “Hello World” Program works.
  - Incorporation into standard still not certain. Need to build case.
  - Next Step: Demonstrate usage with threaded triangular solve.
- Exascale potential:
  - Incremental path to MPI+X.
  - Dial-able SMP scope.

```c
int n = ...;
double* values;
MPI_Comm_alloc_mem(
    MPI_COMM_NODE, // comm (SOCKET works too)
    n * sizeof(double), // size in bytes
    MPI_INFO_NULL, // placeholder for now
    &values); // Pointer to shared array (out)

// At this point:
// - All ranks on a node/socket have pointer to a shared buffer (values).
// - Can continue in MPI mode (using shared memory algorithms) or
// - Can quiet all but one:
int rank;
MPI_Comm_rank(MPI_COMM_NODE, &rank);
if (rank == 0) {
    // Start threaded code segment, only on rank 0 of the node
    ...
}
MPI_Comm_free_mem(MPI_COMM_NODE, values);
```

Collaborators: B. Barrett, Brightwell, Wolf - SNL; Vallee, Koenig - ORNL
Preconditioners for Scalable Multicore Systems

- Observe: Iteration count increases with number of subdomains.
- With scalable threaded smoothers (LU, ILU, Gauss-Seidel):
  - Solve with fewer, larger subdomains.
  - Better kernel scaling (threads vs. MPI processes).
  - Better convergence, More robust.
- Exascale Potential: Tiled, pipelined implementation.
- Three efforts:
  - Level-scheduled triangular sweeps (ILU solve, Gauss-Seidel).
  - Decomposition by partitioning
  - Multithreaded direct factorization

---

Placement and Migration
Placement and Migration

• MPI:
  – Data/work placement clear.
  – Migration explicit.

• Threading:
  – It’s a mess (IMHO).
  – Some platforms good.
  – Many not.
  – Default is bad (but getting better).
  – Some issues are intrinsic.
Data Placement on NUMA

• Memory Intensive computations: Page placement has huge impact.
• Most systems: First touch (except LWKs).
• Application data objects:
  – Phase 1: Construction phase, e.g., finite element assembly.
  – Phase 2: Use phase, e.g., linear solve.
• Problem: First touch difficult to control in phase 1.
• Idea: Page migration.
Data placement experiments

- MiniApp: HPCCG (Mantevo Project)
- Construct sparse linear system, solve with CG.
- Two modes:
  - Data placed by assembly, not migrated for NUMA
  - Data migrated using parallel access pattern of CG.
- Results on dual socket quad-core Nehalem system.
Weak Scaling Problem

- MPI and conditioned data approach comparable.
- Non-conditioned very poor scaling.
Page Placement summary

• MPI+OpenMP (or any threading approach) is best overall.
• But:
  – Data placement is big issue.
  – Hard to control.
  – Insufficient runtime support.
• Current work:
  – Migrate on next-touch (MONT).
  – Considered in OpenMP (next version).
  – Also being studied in Kitten (Kevin Pedretti).
• Note: This phenomenon especially damaging to OpenMP common usage.
Resilient Algorithms
My Luxury in Life (wrt FT/Resilience)

The privilege to think of a computer as a **reliable, digital** machine.

“At 8 nm process technology, it will be harder to tell a 1 from a 0.”

(W. Camp)
Users’ View of the System Now

• “All nodes up and running.”
• Certainly nodes fail, but invisible to user.
• No need for me to be concerned.
• Someone else’s problem.
• Nodes in one of four states.
  1. Dead.
  2. Dying (perhaps producing faulty results).
  3. Reviving.
  4. Running properly:
     a) Fully reliable or…
     b) Maybe still producing an occasional bad result.
Hard Error Futures

• C/R will continue as dominant approach:
  – Global state to global file system OK for small systems.
  – Large systems: State control will be localized, use SSD.

• Checkpoint-less restart:
  – Requires full vertical HW/SW stack co-operation.
  – Very challenging.
  – Stratified research efforts not effective.
Soft Error Futures

- Soft error handling: A legitimate algorithms issue.
- Programming model, runtime environment play role.
Consider GMRES as an example of how soft errors affect correctness

- **Basic Steps**
  1) Compute Krylov subspace (preconditioned sparse matrix-vector multiplies)
  2) Compute orthonormal basis for Krylov subspace (matrix factorization)
  3) Compute vector yielding minimum residual in subspace (linear least squares)
  4) Map to next iterate in the full space
  5) Repeat until residual is sufficiently small

- More examples in Bronevetsky & Supinski, 2008
Why GMRES?

• Many apps are implicit.
• Most popular (nonsymmetric) linear solver is preconditioned GMRES.
• Only small subset of calculations need to be reliable.
  – GMRES is iterative, but also direct.
Every calculation matters

- Small PDE Problem: ILUT/GMRES
- Correct result: 35 Iters, 343M FLOPS
- 2 examples of a single bad op.
- **Solvers:**
  - 50-90% of total app operations.
  - Soft errors most likely in solver.
- **Need new algorithms for soft errors:**
  - Well-conditioned wrt errors.
  - Decay proportional to number of errors.
  - Minimal impact when no errors.

<table>
<thead>
<tr>
<th>Description</th>
<th>Iters</th>
<th>FLOPS</th>
<th>Recursive Residual Error</th>
<th>Solution Error</th>
</tr>
</thead>
<tbody>
<tr>
<td>All Correct Calcs</td>
<td>35</td>
<td>343M</td>
<td>4.6e-15</td>
<td>1.0e-6</td>
</tr>
<tr>
<td>Iter=2, y[1] += 1.0 SpMV incorrect Ortho subspace</td>
<td>35</td>
<td>343M</td>
<td>6.7e-15</td>
<td>3.7e+3</td>
</tr>
<tr>
<td>Q[1][1] += 1.0 Non-ortho subspace</td>
<td>N/C</td>
<td>N/A</td>
<td>7.7e-02</td>
<td>5.9e+5</td>
</tr>
</tbody>
</table>

- New Programming Model Elements: SW-enabled, highly reliable:
  - Data storage, paths.
  - Compute regions.
- Idea: *New algorithms with minimal usage of high reliability.*
- First new algorithm: Flexible-operator (FO)-GMRES.
  - Resilient to soft errors.
  - Only orthogonalization vectors and computations highly reliable.
  - Vast majority of data, ops done with base reliability:
    - Operator, preconditioner data
    - SpMV, Preconditioner application

M. Heroux, M. Hoemmen
Algorithms and Meta-Algorithms
Communication-avoiding iterative methods

- Iterative Solvers:
  - Dominant cost of many apps (up to 80+% of runtime).

- Exascale challenges for iterative solvers:
  - Collectives, synchronization.
  - Memory latency/BW.
  - Not viable on exascale systems in present forms.

- Communication-avoiding (s-step) iterative solvers:
  - Idea: Perform s steps in bulk (s=5 or more):
    - s times fewer synchronizations.
    - s times fewer data transfers: Better latency/BW.
  - Problem: Numerical accuracy of orthogonalization.

- New orthogonalization algorithm:
  - Tall Skinny QR factorization (TSQR).
  - Communicates less and more accurate than previous approaches.
  - Enables reliable, efficient s-step methods.

- TSQR Implementation:
  - 2-level parallelism (Inter and intra node).
  - Memory hierarchy optimizations.
  - Flexible node-level scheduling via Intel Threading Building Blocks.
  - Generic scalar data type: supports mixed and extended precision.

TSQR capability:
- Critical for exascale solvers.
- Part of the Trilinos scalable multicore capabilities.
- Helps all iterative solvers in Trilinos (available to external libraries, too).
- Staffing: Mark Hoemmen (lead, post-doc, UC-Berkeley), M. Heroux
- Part of Trilinos 10.6 release, Sep 2010.
Advanced Modeling and Simulation Capabilities: Stability, Uncertainty and Optimization

- Promise: 10-1000 times increase in parallelism (or more).

SPDEs:

- Pre-requisite: High-fidelity “forward” solve:
  - Computing families of solutions to similar problems.
  - Differences in results must be meaningful.

- Size of a single forward problem
## Advanced Capabilities: Readiness and Importance

<table>
<thead>
<tr>
<th>Modeling Area</th>
<th>Sufficient Fidelity?</th>
<th>Other concerns</th>
<th>Advanced capabilities priority</th>
</tr>
</thead>
<tbody>
<tr>
<td>Seismic</td>
<td>Yes.</td>
<td>None as big.</td>
<td>Top.</td>
</tr>
<tr>
<td><em>S. Collis, C. Ober</em></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Shock &amp; Multiphysics (Alegra)</td>
<td>Yes, but some concerns.</td>
<td>Constitutive models, material responses maturity.</td>
<td>Secondary now. Non-intrusive most attractive.</td>
</tr>
<tr>
<td><em>A. Robinson, C. Ober</em></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Multiphysics (Charon)</td>
<td>Reacting flow w/ simple transport, device w/ drift diffusion, …</td>
<td>Higher fidelity, more accurate multiphysics.</td>
<td>Emerging, not top.</td>
</tr>
<tr>
<td><em>J. Shadid</em></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><em>K. Pierson</em></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Advanced Capabilities: Other issues

• Non-intrusive algorithms (e.g., Dakota):
  – Task level parallel:
    • A true peta/exa scale problem?
    • Needs a cluster of 1000 tera/peta scale nodes.

• Embedded/intrusive algorithms (e.g., Trilinos):
  – Cost of code refactoring:
    • Non-linear application becomes “subroutine”.
    • Disruptive, pervasive design changes.

• Forward problem fidelity:
  – Not uniformly available.
  – Smoothness issues.
  – Material responses.
Advanced Capabilities: Derived Requirements

• Large-scale problem presents collections of related subproblems with forward problem sizes.

• Linear Solvers: \( Ax = b \rightarrow AX = B, \ Ax^i = b^i, \ A^i x^i = b^i \)
  – Krylov methods for multiple RHS, related systems.

• Preconditioners:
  – Preconditioners for related systems.

\[ A^i = A_0 + \Delta A^i \]

• Data structures/communication: \( \text{pattern}(A^i) = \text{pattern}(A^j) \)
  – Substantial graph data reuse.
Summary

- Kernel sets will expand:
  - Advanced modeling and simulation: Gives a better answer.
  - “Communication-avoiding” methods, e.g., matrix-powers kernel, replication.
- Bi-modal MPI/MPI+threading can allow dial-able SMP execution.
- Resilient algorithms will mitigate soft error impact.
- Building the next generation of parallel applications requires enabling domain scientists:
  - Write sophisticated methods.
  - Do so with serial fragments.
  - Fragments hoisted into scalable, resilient fragment.
- Success of manycore will require breaking out of global SIMT-only.
- Migration of Fortran apps to manycore will be painful.
Quiz (True or False)

1. MPI-only has the best parallel performance.
2. Future parallel applications will not have MPI_Init().
3. All future programmers will need to write parallel code.
4. Use of “markup”, e.g., OpenMP pragmas, is the least intrusive approach to parallelizing a code.
5. DRY is not possible across CPUs and GPUs.
6. Extended precision is too expensive to be useful.
7. Resilience will be built into algorithms.
8. GPUs are a harbinger of CPU things to come.
9. Fortran Developers are in trouble in a manycore world.
10. Global SIMT is sufficient parallelism for scientific computing.